

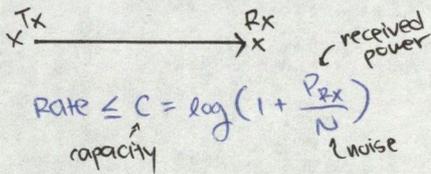
Lecture 3 Notes

Today:

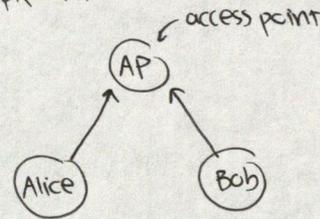
- Broadcast nature of wireless channel \rightarrow interference when you have multiple Tx
- rate region
- interference cancellation (IC)
- Zigzag

So far we've talked about single channel.

single channel:



Multiple Tx:



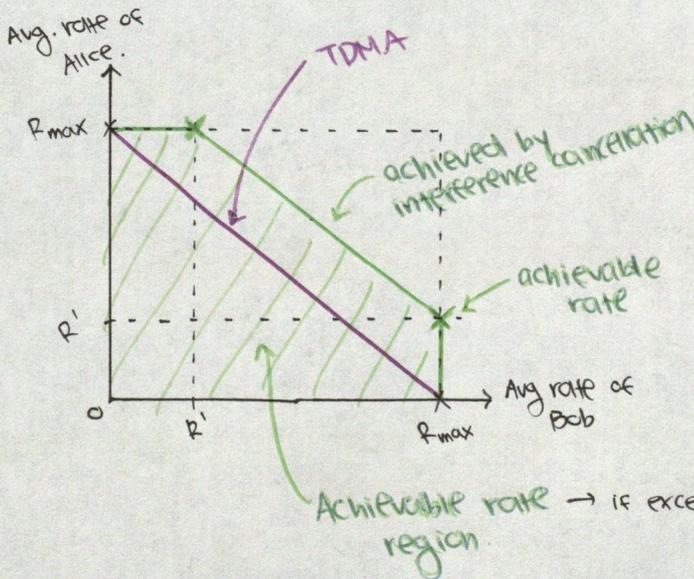
- limited by interference

- assume:

$$P_{Rx}^{Alice} = P_{Rx}^{Bob} = P$$

- can Alice & Bob transmit @ C?
 - \hookrightarrow if Bob is alone, can transmit @ max rate.
 - \hookrightarrow cannot both transmit @ max (would result in unbounded C.)
- rate of channel must be smaller than capacity

Rate Region:



$$R_{max} = \log(1 + \frac{P}{N})$$

- if Alice transmits @ R_{max} :

$$R_{Bob} \leq \log(1 + \frac{P_B}{N + P_A}) = R'$$

(for Bob, A's signal is noise)

* any segment between 2 achievable points is achievable

Achievable rate region \rightarrow if exceeded, either A or B is undecodable.

Remark 1:

① Assume Alice is noise and decode Bob.

$$y(t) = \underbrace{y_A(t)}_{\text{signal from Alice}} + \underbrace{y_B(t)}_{\text{from Bob}} + \underbrace{w(t)}_{\text{noise}}$$

* if decoding Bob, assume $y_A(t)$ & $w(t)$ are same term (both noise): $w'(t) = y_A(t) + w(t)$

\hookrightarrow get $\hat{x}_B(t)$ once decoded.

② subtract Bob's signal (already decoded) $\rightarrow y_A(t) + w(t) \rightarrow$ decode Alice. (typically easier to decode bigger power first)

Remark 2:

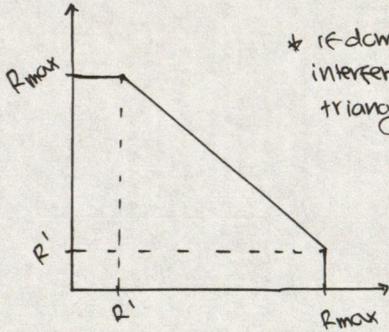
on rate region graph, purple line is achieved by TDMA or CSMA.

↳ TDMA/CSMA → alternating between 2 transmitters.

Remark 3:

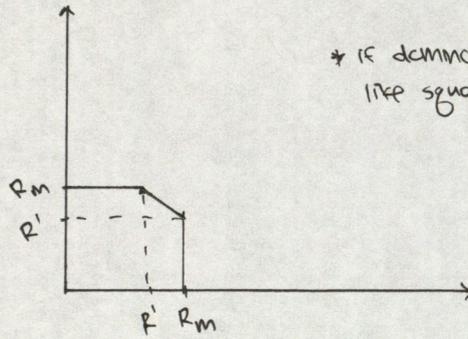
$N \ll P_{rx}$
 $R' \ll R_{max}$

WiFi/LTE



* is dominated by interference, graph like triangle

$N \gg P_{rx}$ (bigger or same order)
 $R' = \log(1 + \frac{P}{P+N}) \sim R_{max}$ (P is negligible/close to noise)



* is dominated by noise, graph like square.

Remark 4:

Generalize to more than 2 users.

- Ex: 3 users (Alice, Bob, Chris)

$R'_{Chris} \leq \log(1 + \frac{P}{2P+N})$

$R'_{Bob} \leq \log(1 + \frac{P}{P+N})$

$R'_{Alice} \leq \log(1 + \frac{P}{N})$

Remark 5:

Interference cancellation needs coordination.

- TDMA needs to coordinate on when to transmit (not at what rate)

- coordinating across transmitters → done by AP.

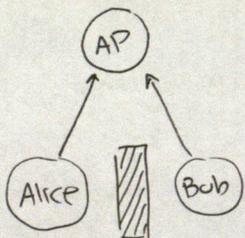
↳ ex: if Alice decides to transmit @ R_{max} , Bob must know, or will exceed rate region → undecodable.

↳ AP/scheme needs to decide who's transmitting first.

Zig Zag:

- tries to perform better but without coordination. (tries to not have AP do something specific)

- Problem: Hidden terminal



- users cannot hear each others' transmissions

↳ Alice & Bob cannot tell when each one transmits in medium

↳ can transmit at same time → collision (packets don't go through)

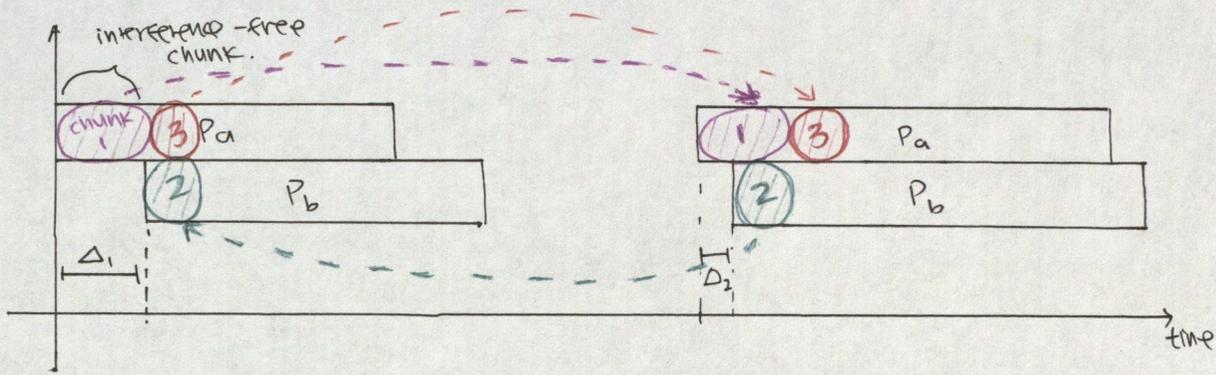
- how to solve HT problem?

↳ if AP were to tell A&B when to transmit → a lot of overhead

↳ must do for every packet for user

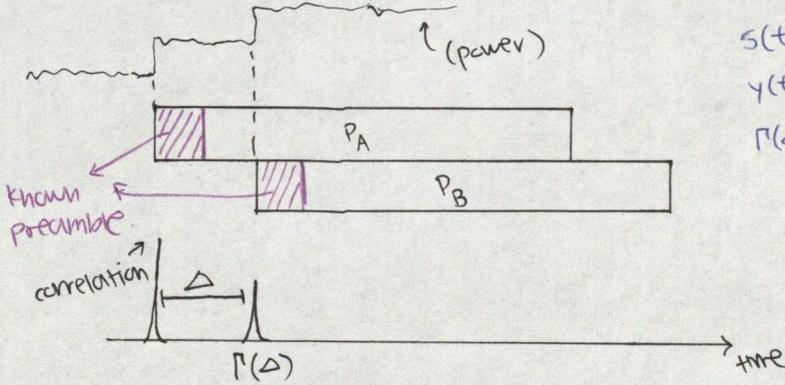
↳ does not know if users still have packets to transfer.

- ZigZag solution: Basic Idea



- will listen before transmitting & randomize transmit time
↳ $\Delta_1 \neq \Delta_2$ thanks to randomization.
- if successful transmission & decode, AP will send acknowledge
↳ if collision → no ack.
- if collision, users send the same data. (same packets)
- can subtract **chunk 1** from 2nd collision (know from first collision)
↳ now know Bob's **chunk 2** and can subtract from collision 1.
↳ find Alice's **chunk 3** and subtract from second collision.
↳ continue until finish decoding

- how do we detect collision and Δ ?



$$s(t) = s_1, \dots, s_k \quad \leftarrow \text{preamble.}$$

$$y(t) = y_A(t) - y_B(t)$$

$$r(\Delta) = \sum y(t, \Delta) s^*(t) \quad \leftarrow \text{summed over size of } s.$$

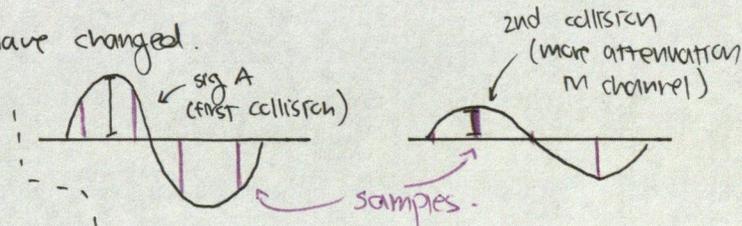
$$= \sum s_B(t) H_B e^{j2\pi\Delta s t} \quad s^*(t) = \sum |s(t)|^2 e^{j2\pi\Delta s t}$$

- correlation w/ preamble.
↳ second spike smaller due to interference.
↳ Δ usually $>$ preamble.

- correlation can be used to determine if 2 collisions are between 2 packets.
↳ correlate each packet w/ each other → if big spike, same packets.

- what does it mean to subtract a chunk from second collision?

- between first & second collision, channel may have changed.
↳ don't sample 2 collisions @ same time.

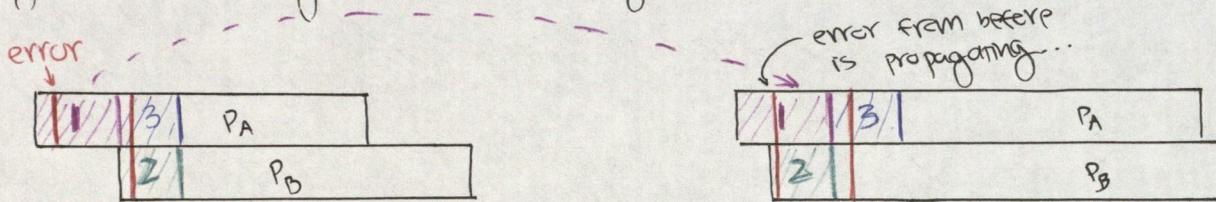


- 1) decode signal (Rx) → bits.
- 2) generate a channel free signal
- 3) apply channel during 2nd collision to channel free signal (pretend sig went through channel @ time of 2nd coll.)
- 4) sample as in 2nd collision.

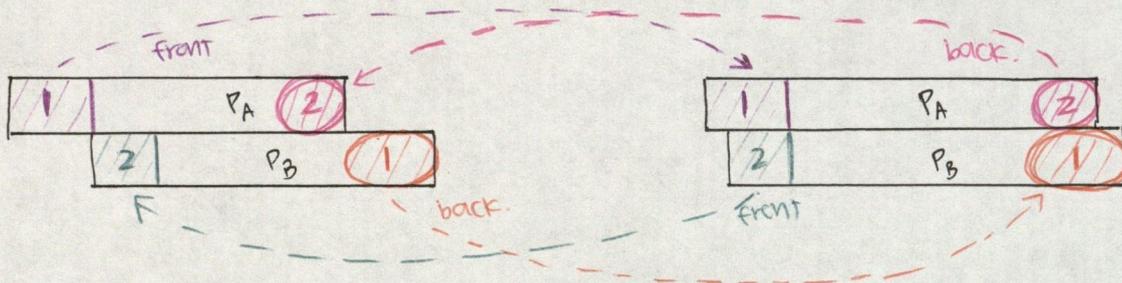
- since sampling may be different, cannot simply subtract sample by sample.

Error Propagations

- Big problem if decoding has error \rightarrow decoding is iterative.



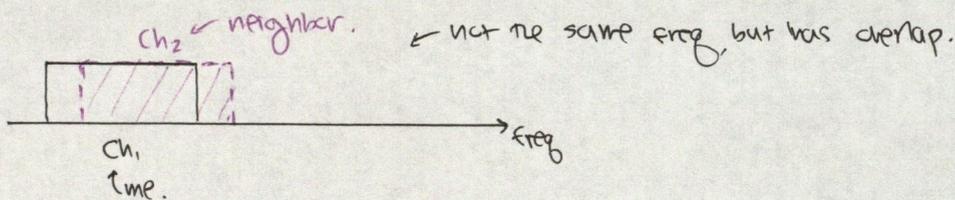
- how to fix this? \Rightarrow also decode from back of packets.



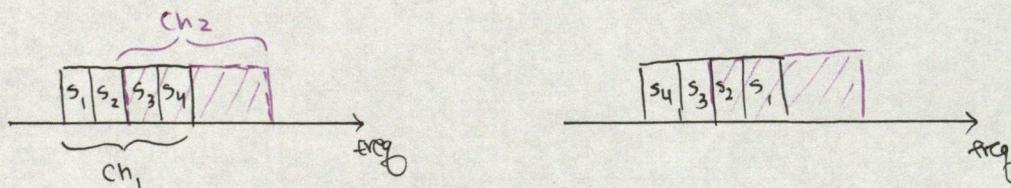
- after decode from front to back, compare.
 \hookrightarrow to increase reliability \rightarrow add preamble to back.

In Freq

- 802.11 has overlapping channels.



- dealing w/ collisions in freq \rightarrow randomize what's transmitted in subcarriers.
 \hookrightarrow no Δ randomization ($\Delta_1 = \Delta_2$), since same channel.



- switch which bits collide with which.