

Consistent Correspondence

Justin Solomon

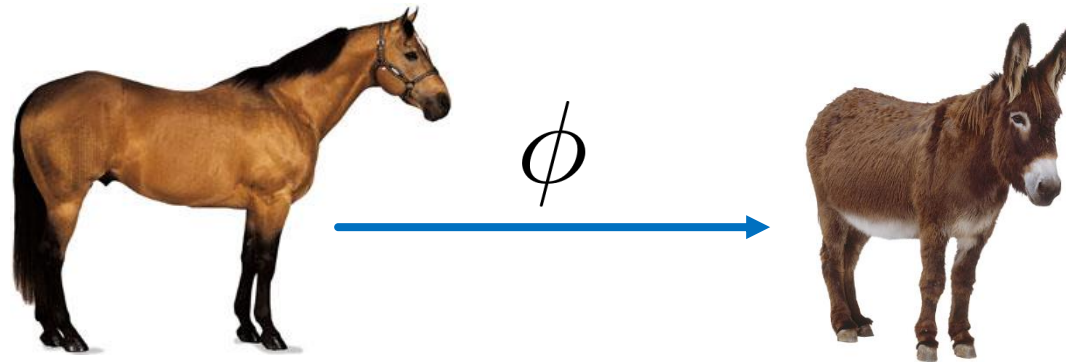
6.8410: Shape Analysis

Spring 2023



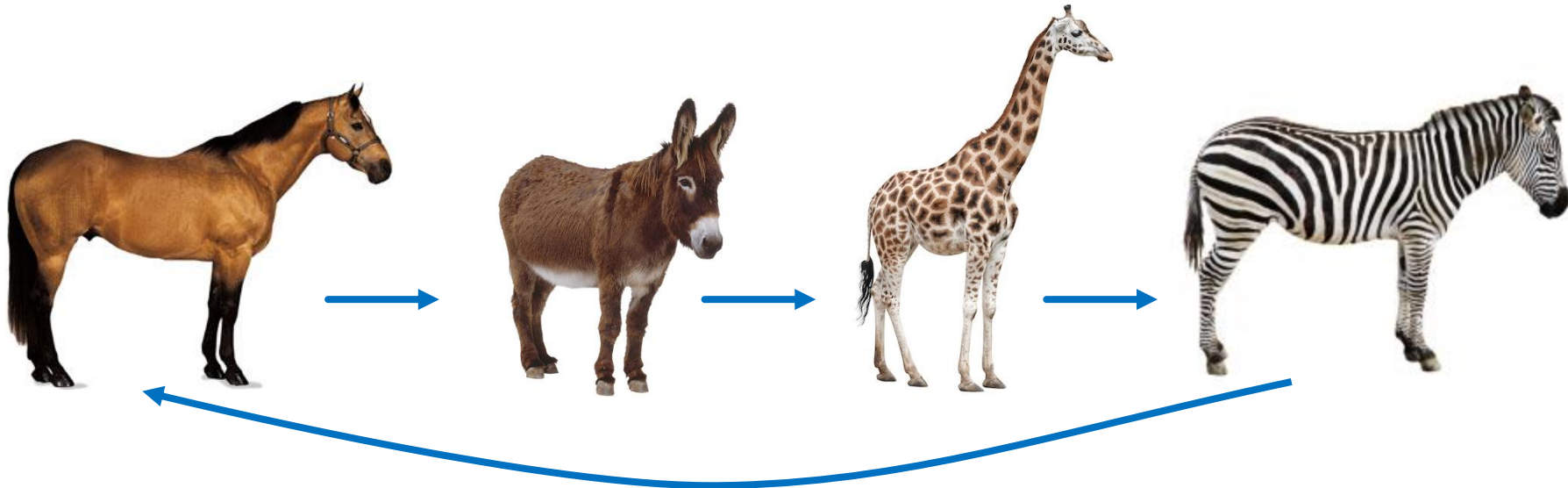
Previously

Map between **two** shapes.



Question

What happens if you **compose** these maps?





What do you **expect** if
you compose
around a cycle?

Cycle consistency

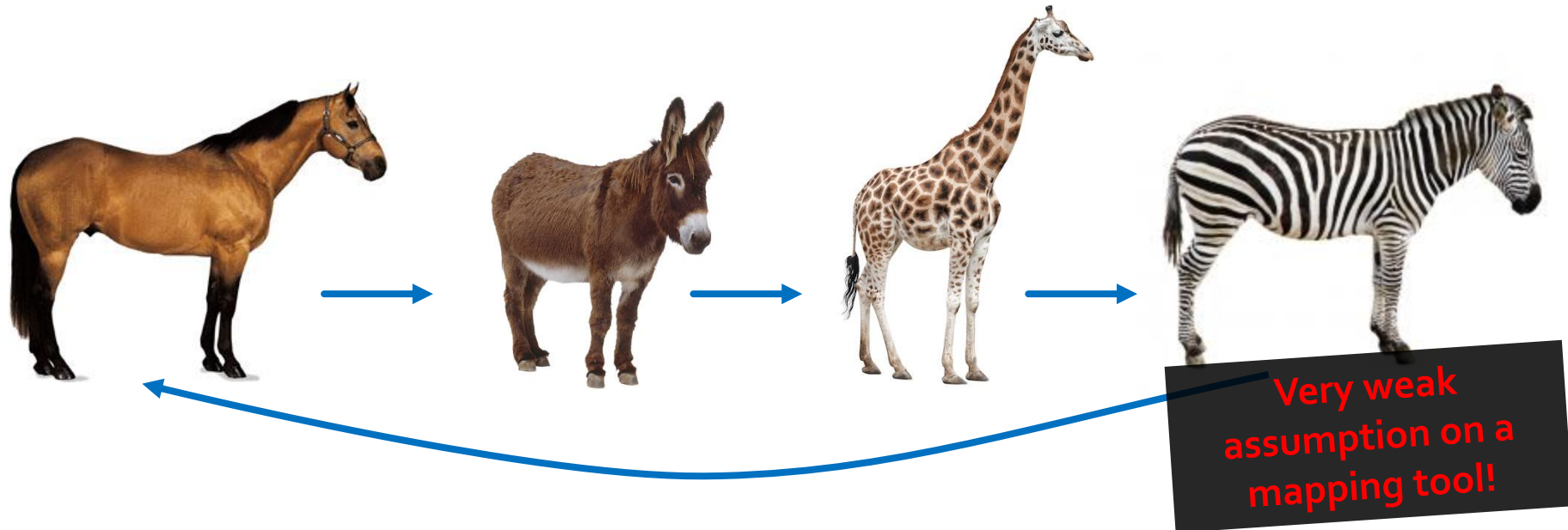
[*sahy-kuh | kuh n-sis-tuh n-see*]:

Composing maps in a cycle
yields the identity



Philosophical Point

You should have a good reason if your correspondences are **inconsistent**.



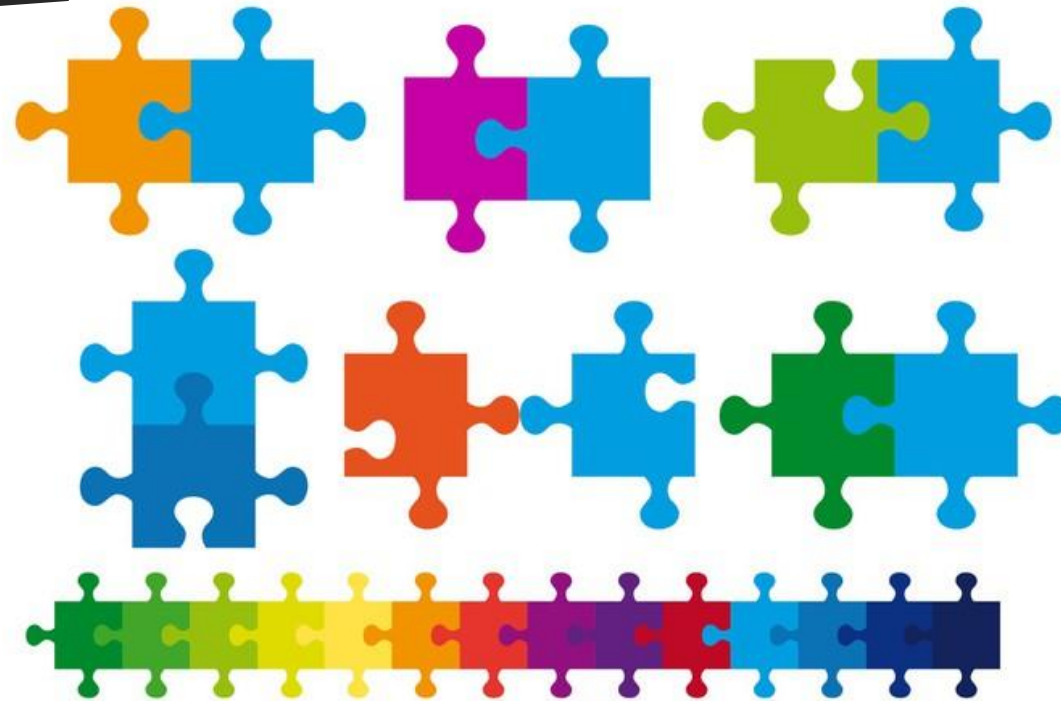
An Unpleasant Constraint

$$\phi_1(\phi_2(\phi_3(x))) = \text{Id}$$

Cycle consistency

Contrasting Viewpoint

Many possible
pairwise matches!



https://s3.pixers.pics/pixers/700/FO/39/51/09/46/700_FO39510946_cd54b90a83d46f5dbd96440271eadfec.jpg

Additional data should help!

Today

Sampling of methods for consistent correspondence.

- Spanning tree
- Inconsistent cycle detection
- Convex optimization

Holy Grail

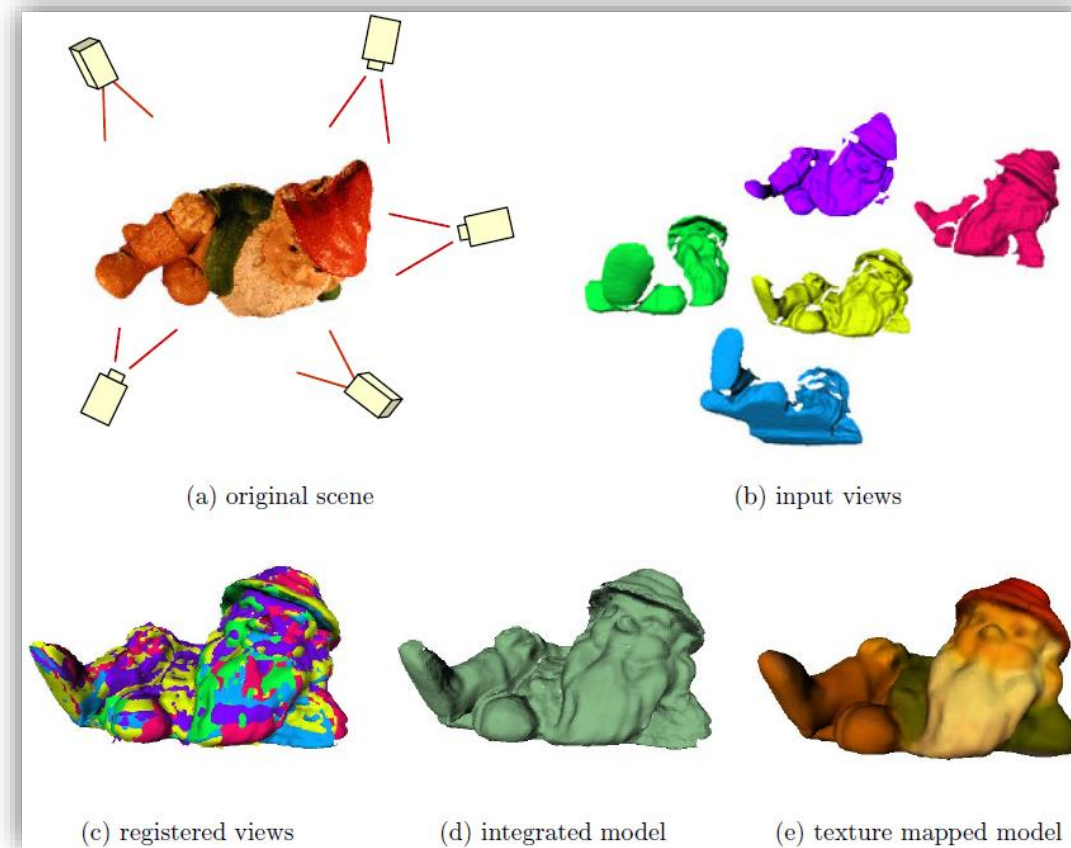
**Simultaneously optimize
all maps in a collection.**

Open problem!

Joint Matching: Simplest Formulation

- **Input**
 - N shapes
 - N^2 maps (see last lecture)
- **Output**
 - Cycle-consistent approximation

Spanning Tree: Original Context



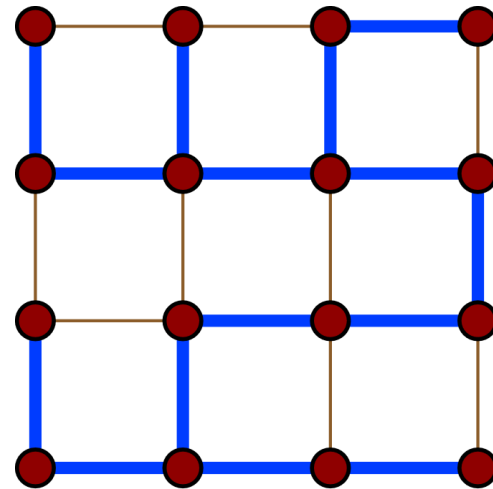
“Automatic Three-Dimensional Modeling from Reality” (Huber, 2002)

Multi-view registration

Unsurprisingly...

Given: Model graph $G = (S, E)$

Find: Largest consistent spanning tree

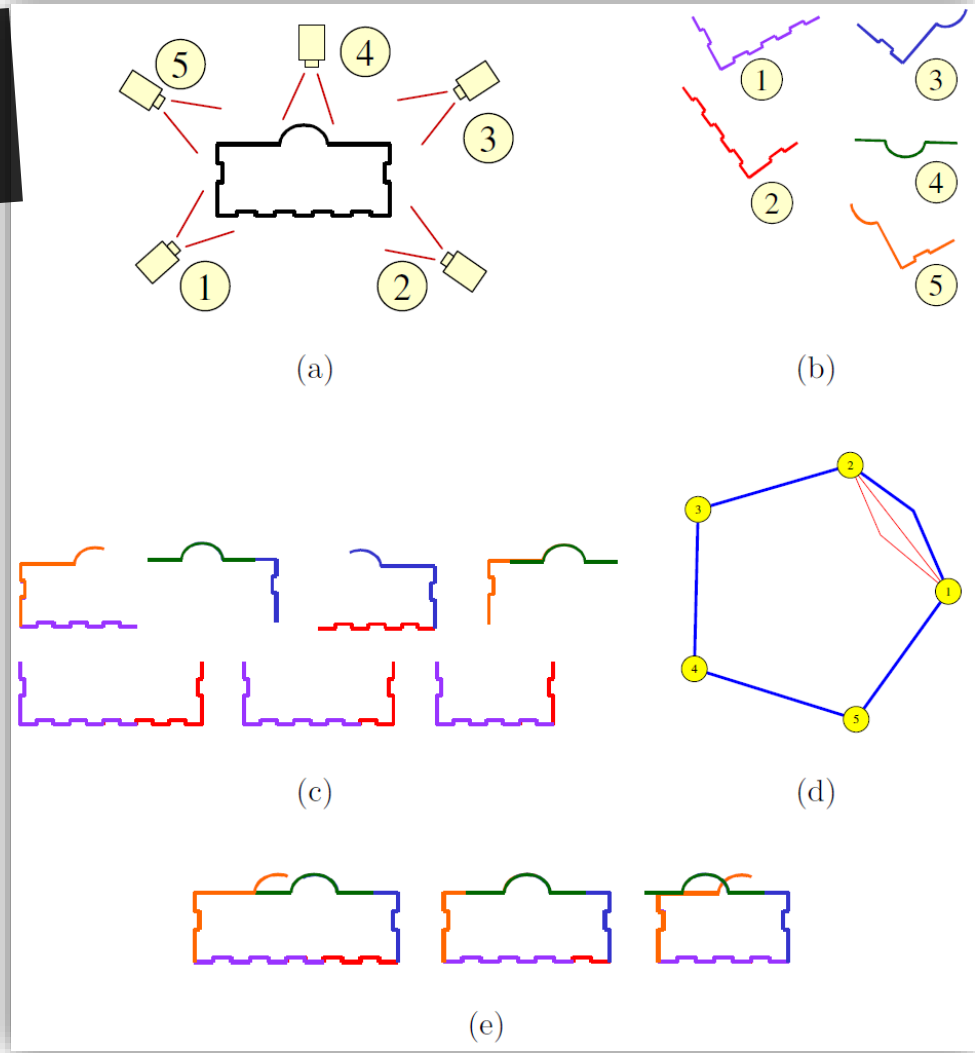


“Automatic Three-Dimensional Modeling from Reality” (Huber, 2002)

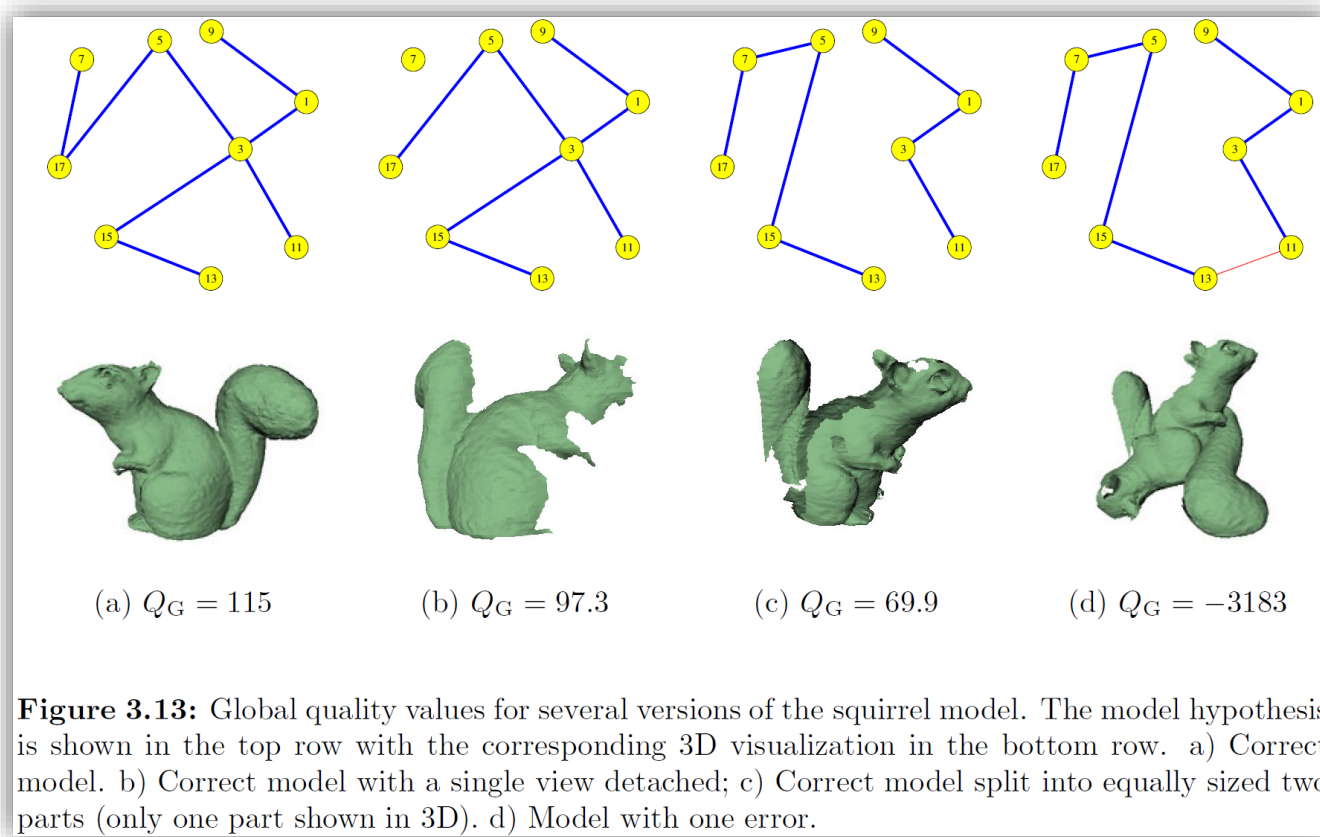
NP-hard

Heuristic Algorithm

Extract consistent spanning tree in model graph



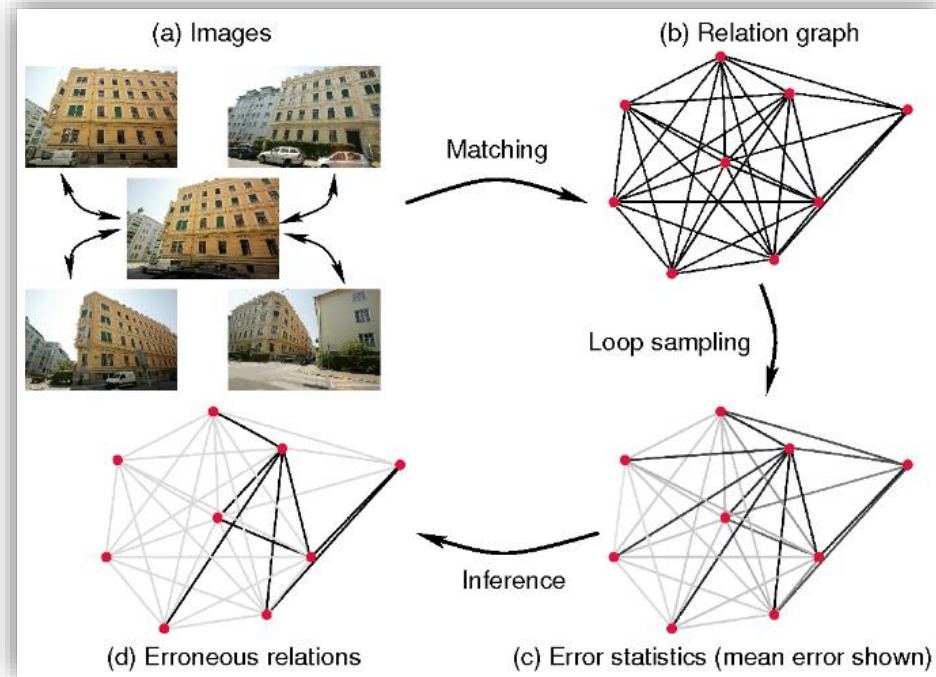
Issues



- Many spanning trees
- **Single incorrect match** can destroy the maps

Inconsistent Loop Detection

Used to deal with repeating structures like windows!



Large for inconsistent cycles

$$\begin{aligned} \max \quad & \sum_L \rho_L x_L \\ \text{s.t.} \quad & x_L \geq x_e \quad \forall e \in L \\ & x_L \leq \sum_{e \in L} x_e \\ & x_L, x_e \in [0, 1] \end{aligned}$$

$x_e = 1$ for false positive edge

$x_L = \max$ of x_e over loop

Relationship: Consistency vs. Accuracy

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An Optimization Approach to Improving Collections of Shape Maps

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Abstract

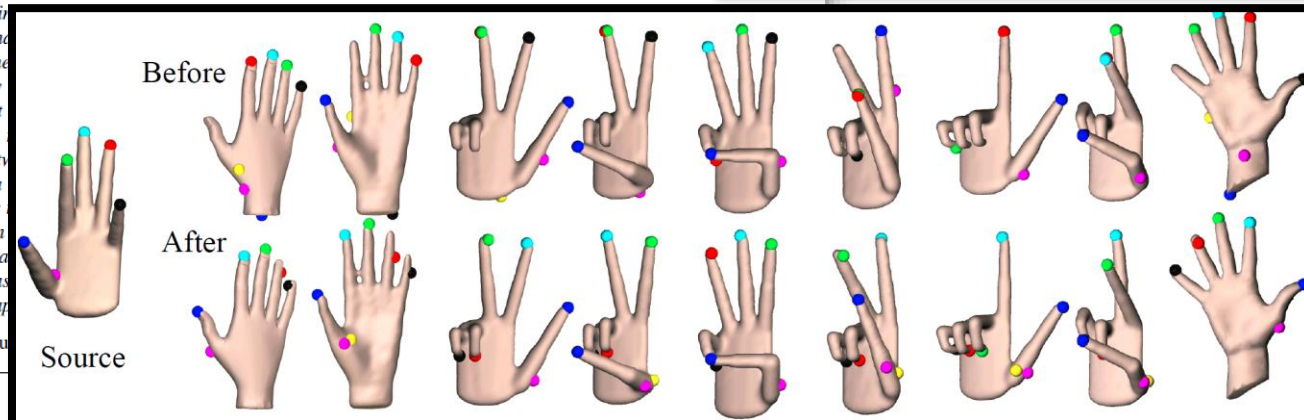
Finding an informative, structure-preserving map between two shapes has been a long-standing problem in geometry processing, involving a variety of solution approaches and applications. However, in many cases, given not only two related shapes, but a collection of them, and considering each pairwise map independently, do not take full advantage of all existing information. For example, a notorious problem with computing shape maps is the ambiguity in

there exist two maps between them based on the sensitivity to how the map consistency is chosen in the new sense interpolate the optimization problem and individually at shapes, as long as for improving map

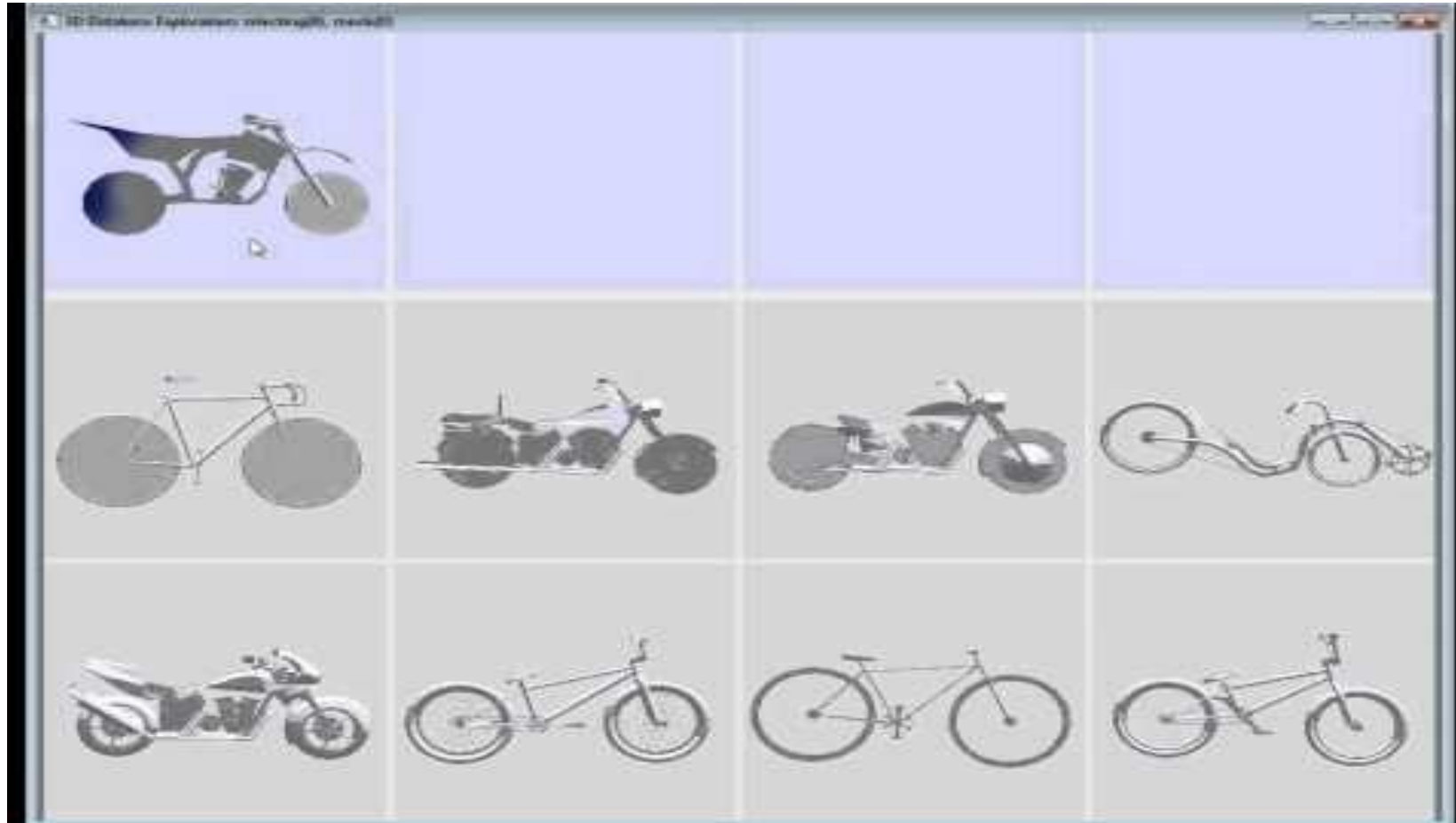
Categories and Su

Iteratively fix triplets and reweight

Definition 3 Given a collection of maps \mathcal{M} , let $\mathcal{B}(\mathcal{M}) = \{m_{i,j} \in \mathcal{M} \mid E_{acc}(m_{i,j}) > 0\}$ — the collection of inaccurate maps. Then we say that \mathcal{M} is *almost accurate*, if there do not exist two maps $m_1, m_2 \in \mathcal{B}(\mathcal{M})$, which both belong to the same 3-cycle in $G_{\mathcal{M}}$. We call such maps *isolated*.



Fuzzy Correspondences

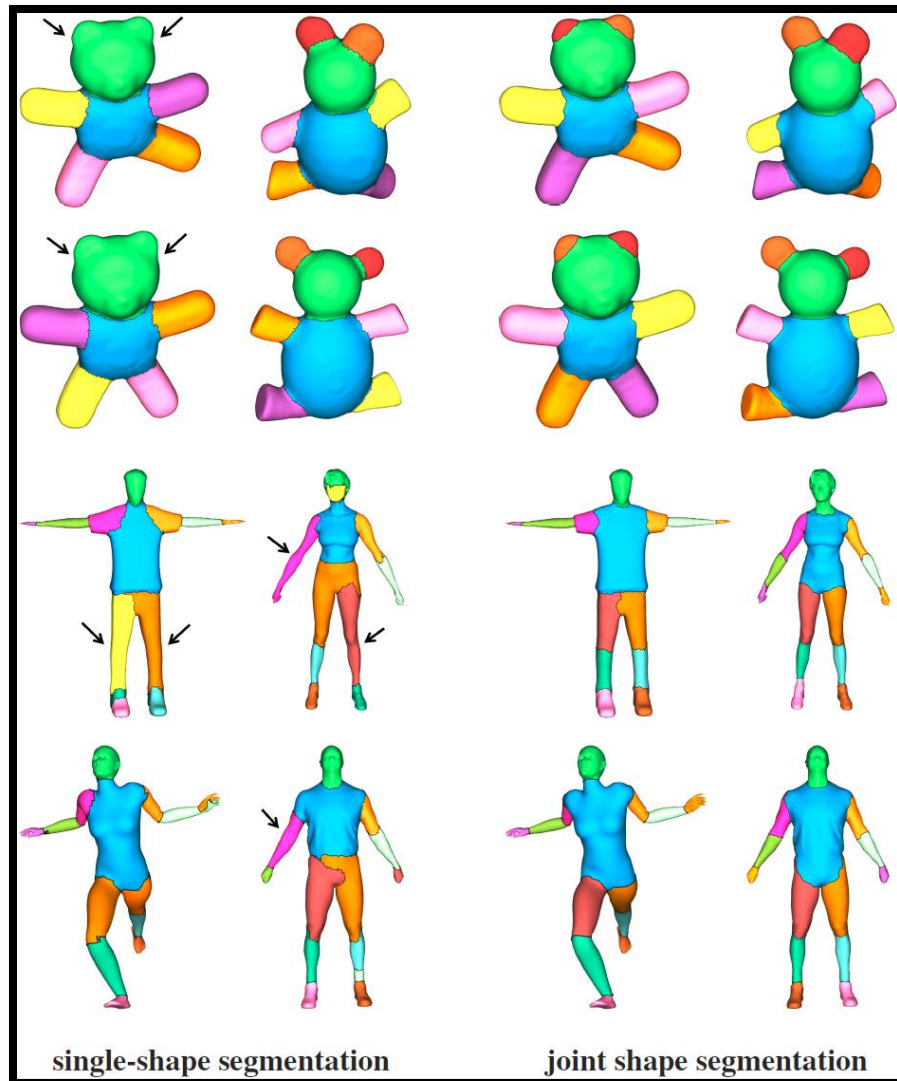


Exploring Collections of 3D Models using Fuzzy Correspondences (Kim et al., SIGGRAPH 2012)

Fuzzy Correspondences: Idea

- Compute $N_k \times N_k$ **similarity matrix**
 - Same number of samples per surface
 - Align similar shapes
- Compute **spectral** embedding
- Use as **descriptor**: Display $e^{-|d_i - d_j|^2}$

Consistent Segmentation



Global optimization
to choose among
many possible
segmentations

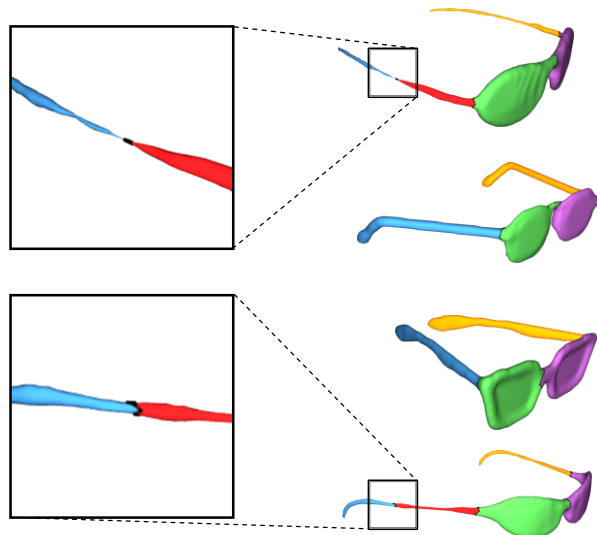
“Joint Shape Segmentation with Linear Programming”
(Huang, Koltun, Guibas; SIGGRAPH Asia 2011)

Joint Segmentation: Motivation

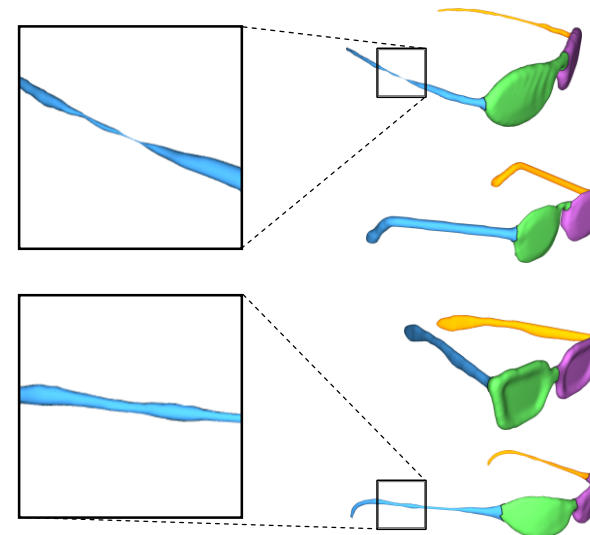
Structural similarity of segmentations

- Extraneous geometric clues

Single shape segmentation
[Chen et al. 09]



Joint shape segmentation
[Huang et al. 11]

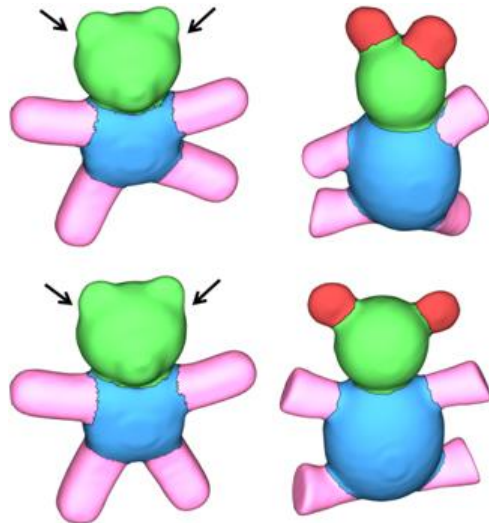


Joint Segmentation: Motivation

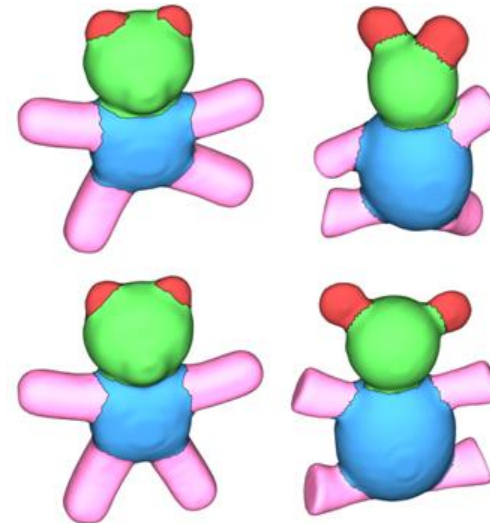
Structural similarity of segmentations

- Low saliency

Single shape segmentation
[Chen et al. 09]



Joint shape segmentation
[Huang et al. 11]

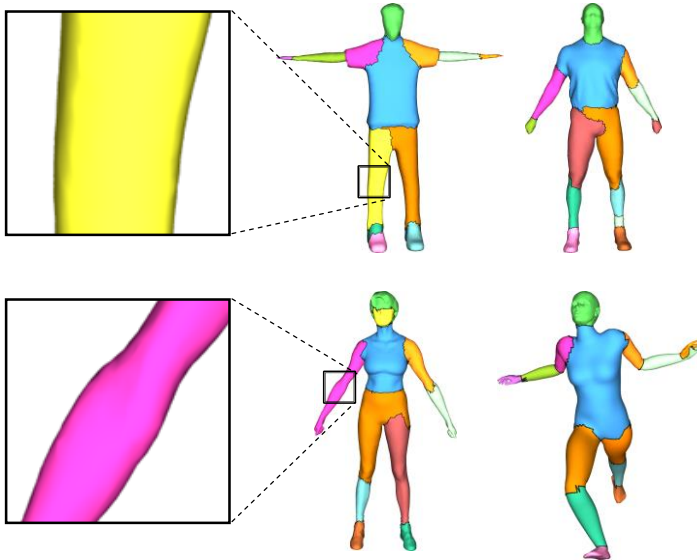


Joint Segmentation: Motivation

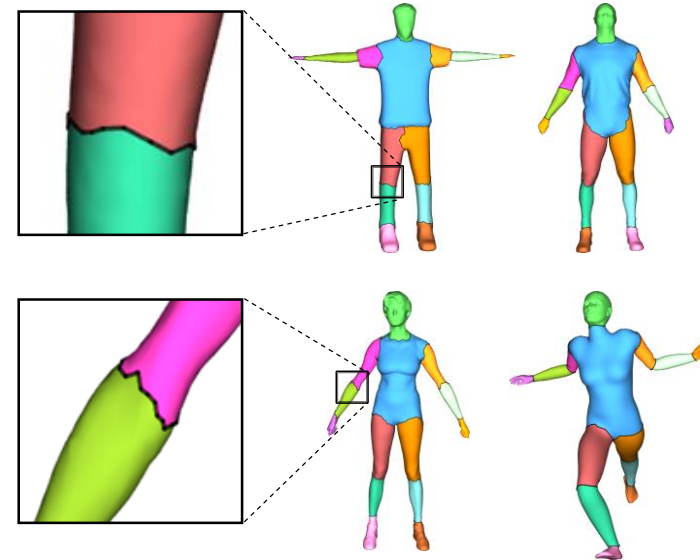
(Rigid) invariance of segments

- Articulated structures

Single shape segmentation
[Chen et al. 09]



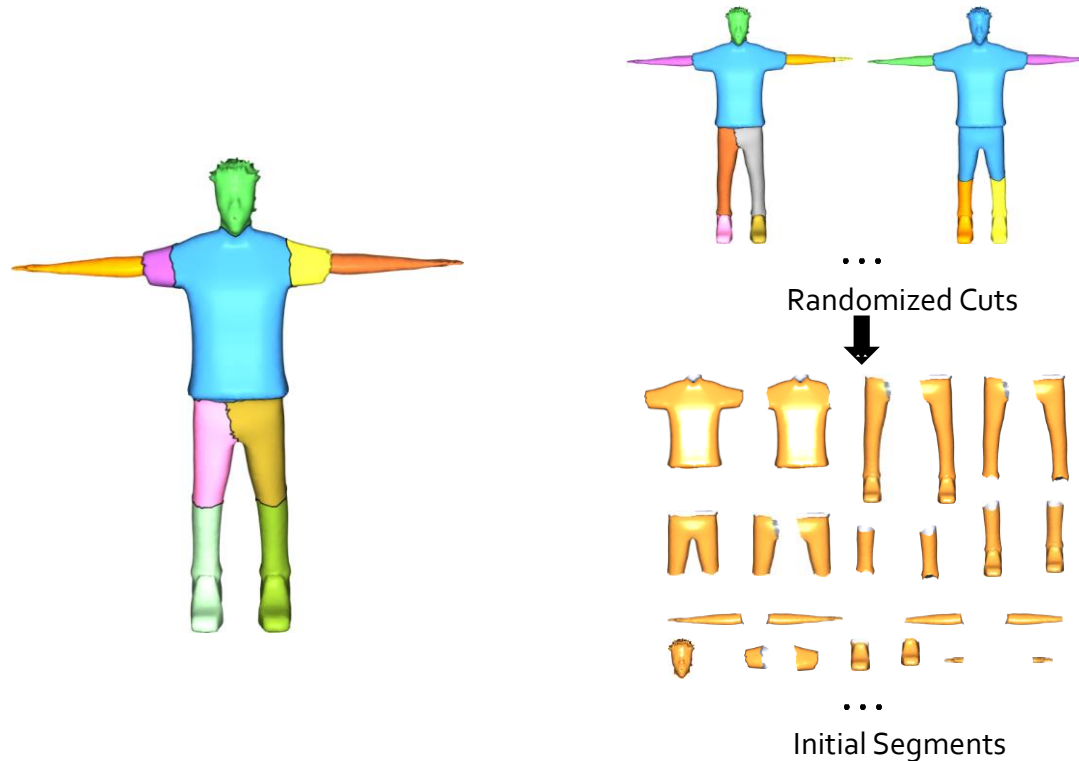
Joint shape segmentation
[Huang et al. 11]



"Joint Shape Segmentation with Linear Programming"
(Huang, Koltun, Guibas; SIGGRAPH Asia 2011; slides provided by authors)

Parameterization

Initial subsets of randomized segmentations



"Joint Shape Segmentation with Linear Programming"
(Huang, Koltun, Guibas; SIGGRAPH Asia 2011; slides provided by authors)

Segmentation Constraint/Score

- Each point covered by one segment

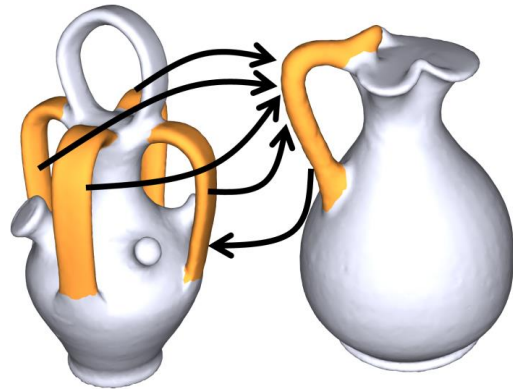
$$|\text{cover}(p)| = 1 \quad \forall p \in W$$

- Avoid tiny segments

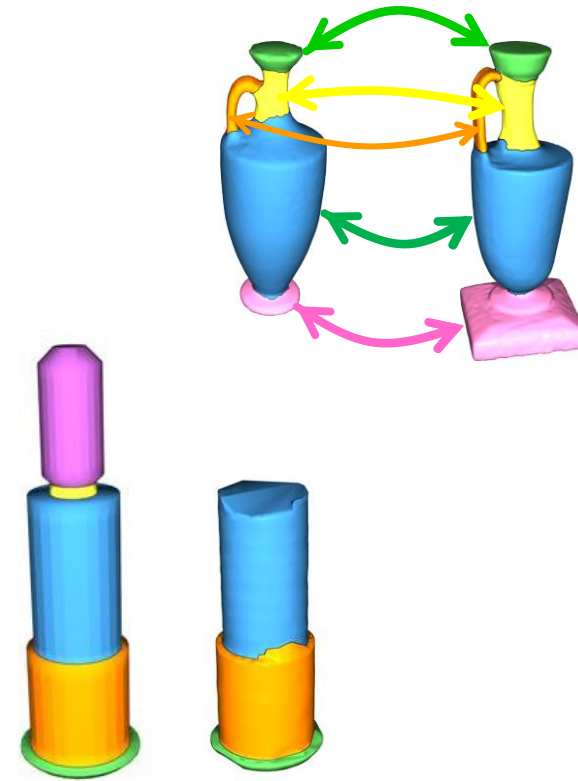
$$\text{score}(S) = \sum_{s \in S} \text{area}(s) \cdot \text{repetitions}_s$$

Consistency Term

- Defined in terms of mappings
 - Oriented
 - Partial



Many-to-one correspondences

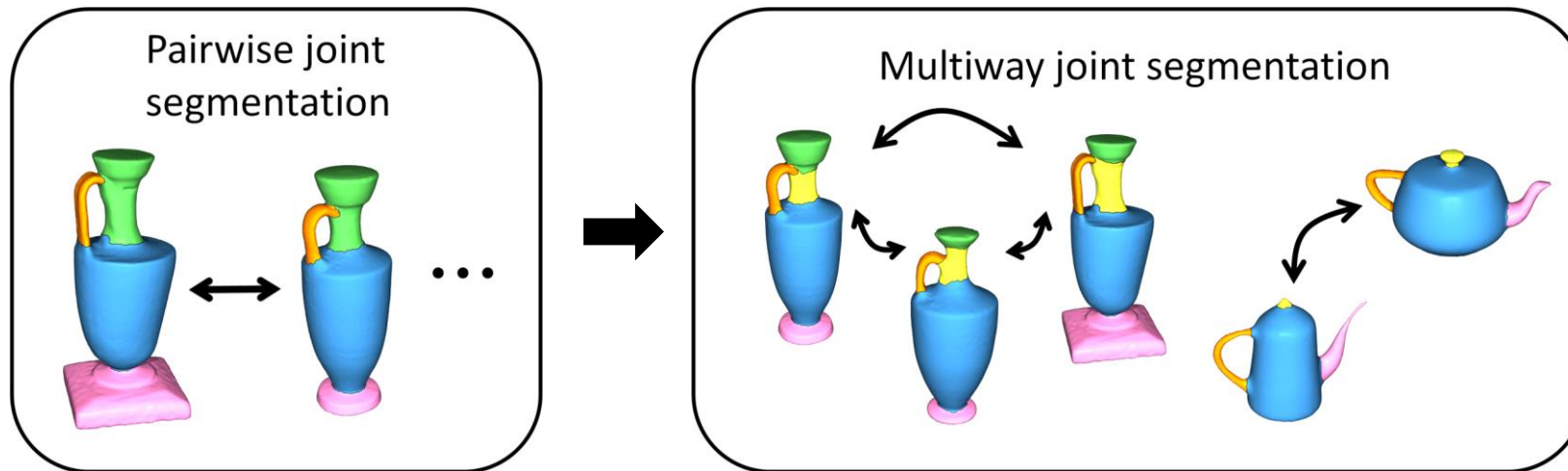


Partial similarity

Multi-Way Joint Segmentation

- Objective function

$$\sum_{i=1}^n \text{score}(S_i) + \sum_{(S_i, S_j) \in \mathcal{E}} \text{consistency}(S_i, S_j)$$



See paper: **Linear program relaxation**



Can you extract
consistent maps in a
globally optimal way?

Basic Setup

$$\begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ 0 & 0 & 0 & \cdots & 1 \\ 0 & 1 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 1 & \cdots & 0 \end{pmatrix}$$

Map as a permutation matrix



What is the **inverse** of a
permutation matrix?

Discrete Relaxation

$$\begin{pmatrix} 1 & 0 & 0 & \dots & 0 \\ 0 & 0 & 0 & \dots & 1 \\ 0 & 1 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 1 & \dots & 0 \end{pmatrix}$$

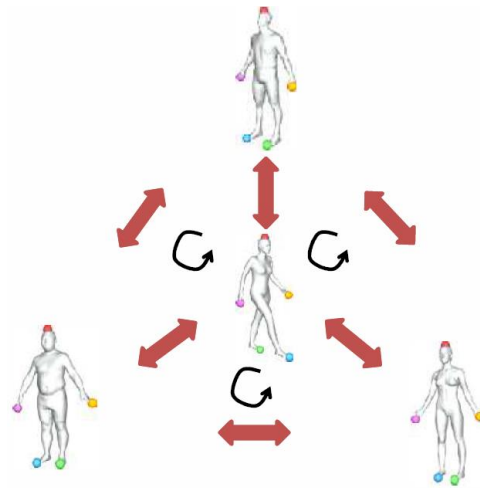
Sums to 1

Sums to 1

Map as a doubly-stochastic matrix

Basic Setting

- Given n objects
- Each object sampled with m points



Map Collection: Matrix Representation

Diagram illustrating the transformation of a map collection from S_1 to S_2 . The transformation is represented by the matrix X_{12} .

$$X_{12} = \begin{bmatrix} 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

$$X = \begin{bmatrix} I_m & X_{12} & \cdots & X_{1n} \\ X_{12}^T & I_m & \cdots & \vdots \\ \vdots & \vdots & \cdots & X_{(n-1),n} \\ X_{1n}^T & \vdots & X_{(n-1),n}^T & I_m \end{bmatrix}$$

- Diagonal blocks are identity matrices
- Off diagonal blocks are permutation matrices
- Symmetric



What is the **rank** of a
consistent map
collection matrix?

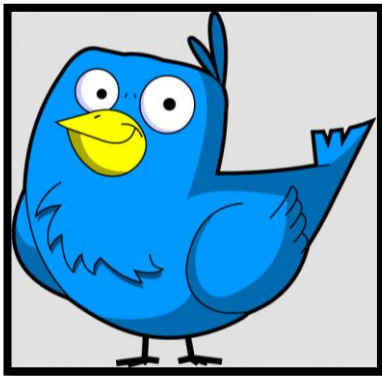
Hint: "Urshape" Factorization

$$X = \begin{bmatrix} I_m & X_{12} & \cdots & X_{1n} \\ X_{12}^T & I_m & \cdots & \vdots \\ \vdots & \vdots & \ddots & X_{(n-1),n} \\ X_{1n}^T & \vdots & X_{(n-1),n}^T & I_m \end{bmatrix}$$

➤ Diagonal blocks are identity matrices

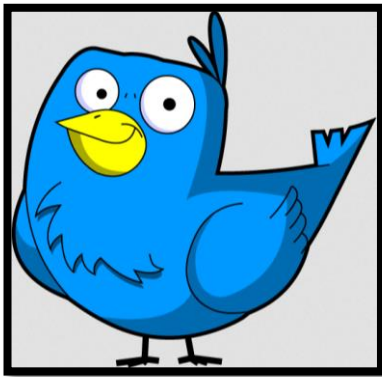
➤ Off diagonal blocks are permutation matrices

➤ Symmetric



Rank m , Number of Samples

$$X_{ij} = X_{j1}^\top X_{i1} \iff X = \begin{pmatrix} I_m \\ \vdots \\ X_{n1}^\top \end{pmatrix} (I_m \cdots X_{n1})$$



Many Equivalent Conditions

Definition 2.1 Given a shape collection $\mathcal{S} = \{S_1, \dots, S_n\}$ of n shapes where each shape consists of the same number of samples, we say a map collection $\Phi = \{\phi_{ij} : S_i \rightarrow S_j \mid 1 \leq i, j \leq n\}$ of maps between all pairs of shapes is cycle consistent if and only if the following equalities are satisfied:

$$\phi_{ii} = id_{S_i}, \quad 1 \leq i \leq n, \quad (1\text{-cycle})$$

$$\phi_{ji} \circ \phi_{ij} = id_{S_i}, \quad 1 \leq i < j \leq n, \quad (2\text{-cycle})$$

$$\phi_{ki} \circ \phi_{jk} \circ \phi_{ij} = id_{S_i}, \quad 1 \leq i < j < k \leq n, \quad (3\text{-cycle}) \quad (1)$$

where id_{S_i} denotes the identity self-map on S_i .

Equivalence for binary map matrix Φ :

1. Φ is cycle-consistent
2. $X = Y_i^\top Y_i$, where $Y_i = (X_{i1}, \dots, X_{in})$
3. $X \succeq 0$

Approximation by Consistent Maps

$$\begin{aligned} \max_X \quad & \sum_{ij \in E} \langle X_{ij}^{\text{in}}, X_{ij} \rangle \\ \text{s.t.} \quad & X \in \{0, 1\}^{nm \times nm} \\ & X \succeq 0 \\ & X_{ii} = I_m \\ & X_{ij} \mathbf{1} = \mathbf{1} \\ & X_{ij}^\top \mathbf{1} = \mathbf{1} \end{aligned}$$

Approximation by Consistent Maps

$$\begin{aligned} \max_X \quad & \sum_{ij \in E} \langle X_{ij}^{\text{in}}, X_{ij} \rangle \\ \text{s.t.} \quad & X \in \{0, 1\}^{nm \times nm} \\ & X \succeq 0 \\ & X_{ii} = I_m \\ & X_{ij} \mathbf{1} = \mathbf{1} \\ & X_{ij}^\top \mathbf{1} = \mathbf{1} \end{aligned}$$

Maximize number
of preserved
matches

Approximation by Consistent Maps

$$\begin{aligned} \max_X & \sum_{ij \in E} \langle X_{ij}^{\text{in}}, X_{ij} \rangle \\ \text{s.t.} & X \in \{0, 1\}^{nm \times nm} \\ & X \succeq 0 \\ & X_{ii} = I_m \\ & X_{ij} \mathbf{1} = \mathbf{1} \\ & X_{ij}^\top \mathbf{1} = \mathbf{1} \end{aligned}$$

Binary
matrix

Approximation by Consistent Maps

$$\begin{aligned} \max_X \quad & \sum_{ij \in E} \langle X_{ij}^{\text{in}}, X_{ij} \rangle \\ \text{s.t.} \quad & X \in \{0, 1\}^{nm \times nm} \\ & X \succeq 0 \\ & X_{ii} = I_m \\ & X_{ij} \mathbf{1} = \mathbf{1} \\ & X_{ij}^\top \mathbf{1} = \mathbf{1} \end{aligned}$$

Every block is a permutation

Approximation by Consistent Maps

$$\begin{aligned} \max_X \quad & \sum_{ij \in E} \langle X_{ij}^{\text{in}}, X_{ij} \rangle \\ \text{s.t.} \quad & X \in \{0, 1\}^{nm \times nm} \\ & X \succeq 0 \\ & X_{ii} = I_m \\ & X_{ij} \mathbf{1} = \mathbf{1} \\ & X_{ij}^\top \mathbf{1} = \mathbf{1} \end{aligned}$$

Self maps
are identity

Approximation by Consistent Maps

$$\begin{aligned} \max_X \quad & \sum_{ij \in E} \langle X_{ij}^{\text{in}}, X_{ij} \rangle \\ \text{s.t.} \quad & X \in \{0, 1\}^{nm \times nm} \end{aligned}$$

$$X \succeq 0$$

Already showed:
Equivalent to low-rank

$$X_{ii} = I_m$$

$$X_{ij} \mathbf{1} = \mathbf{1}$$

$$X_{ij}^\top \mathbf{1} = \mathbf{1}$$

Approximation by Consistent Maps

$$\begin{aligned} \max_X \quad & \sum_{ij \in E} \langle X_{ij}^{\text{in}}, X_{ij} \rangle \\ \text{s.t.} \quad & X \in \{0, 1\}^{nm \times nm} \end{aligned}$$

Nonconvex!

$$X \succeq 0$$

$$X_{ii} = I_m$$

$$X_{ij} \mathbf{1} = \mathbf{1}$$

$$X_{ij}^\top \mathbf{1} = \mathbf{1}$$

Convex Relaxation

$$\begin{aligned} \max_X \quad & \sum_{ij \in E} \langle X_{ij}^{\text{in}}, X_{ij} \rangle \\ \text{s.t.} \quad & X \succeq 0 \\ & X \succcurlyeq 0 \\ & X_{ii} = I_m \\ & X_{ij} \mathbf{1} = \mathbf{1} \\ & X_{ij}^\top \mathbf{1} = \mathbf{1} \end{aligned}$$

Rounding Procedure

Guaranteed to
give permutation

$$\begin{aligned} \max_X \quad & \langle X, X_0 \rangle \\ \text{s.t.} \quad & X \geq 0 \\ & X \mathbf{1} = \mathbf{1} \\ & X^\top \mathbf{1} = \mathbf{1} \end{aligned}$$

Linear assignment problem

Recovery Theorem

Can tolerate $\lambda_2/4(n - 1)$ incorrect correspondences from each sample on one shape.

λ_2 is algebraic connectivity; bounded above by two times maximum degree

`\omit{proof}`

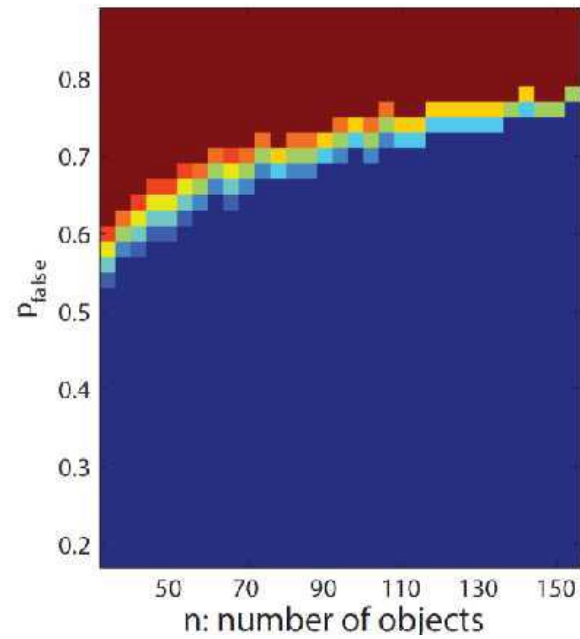
Recovery Theorem: Complete Graph

Can tolerate **25%** incorrect correspondences from each sample on one shape.

λ_2 is algebraic connectivity; bounded above by two times maximum degree

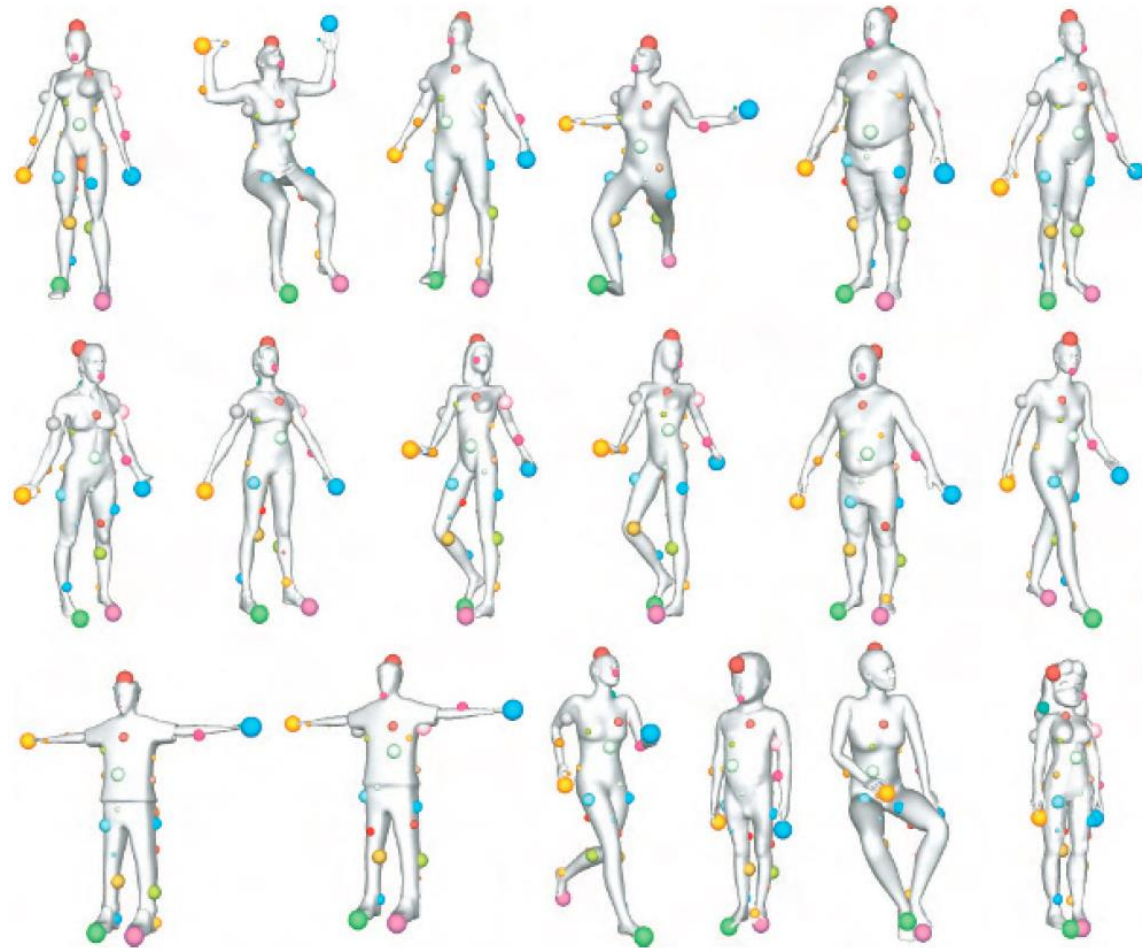
`\omit{proof}`

Phase Transition



Always recovers / **Never recovers**

Example Result



Weaker Relaxation

Solving the multi-way matching problem by permutation synchronization

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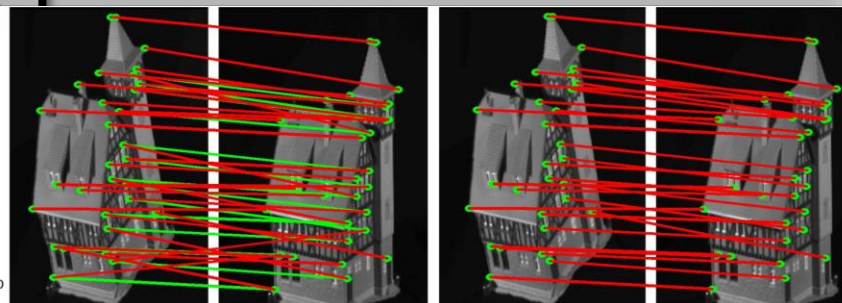
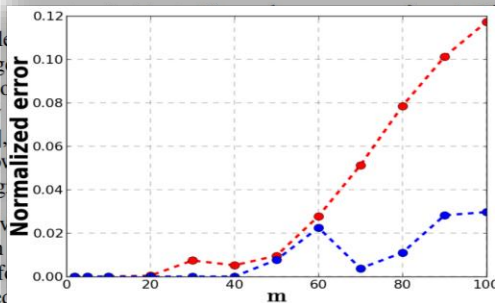
Abstract

The problem of matching not just two, but m different sets of objects to each other arises in many contexts, including finding the correspondence between feature points across multiple images in computer vision. At present it is usually solved by matching the sets pairwise, in series. In contrast, we propose a new method, Permutation Synchronization, which finds all the matchings jointly, in one shot, via a relaxation to eigenvector decomposition. The resulting algorithm is both computationally efficient, and, as we demonstrate with theoretical arguments as well as experimental results, much more stable to noise than previous methods.

1 Introduction

Finding the correct bijection between $\{x'_1, x'_2, \dots, x'_n\}$ is a fundamental problem in many contexts [1]. In this paper, we consider its generalization to m sets of objects X_1, X_2, \dots, X_m . Our primary motivation is the problem of matching landmarks (feature points) across many images in computer vision, a key ingredient of image registration [2], structure from motion (SFM) [8, 9]. However, this problem also arises to problems such as matching multiple graphs [3].

Presently, multi-matching is usually solved by matching X_1 to X_2 , then a permutation synchronization step can be used to find the matchings between the remaining sets. However, if the data are noisy, a single error in the sequential pairwise matches [12, 13, 14]. In contrast, in this paper we describe a new method, Permutation Synchronization, which finds all the matchings jointly, in one shot, via a relaxation to eigenvector decomposition. The resulting algorithm is both computationally efficient, and, as we demonstrate with theoretical arguments as well as experimental results, much more stable to noise than previous methods.



Eigenvector relaxation
of the same problem



Where do the pairwise
input maps come from?

Possible Extension with Guarantees

Eurographics Symposium on Geometry Processing 2015
Mirela Ben-Chen and Ligang Liu
(Guest Editors)

Volume 34 (2015), Number 5

Heavy optimization problem!

Tight Relaxation of Quadratic Matching

Itay Kezurer[†] Shahar Z. Kovalsky[†] Ronen Basri Yaron Lipman

Weizmann Institute of Science

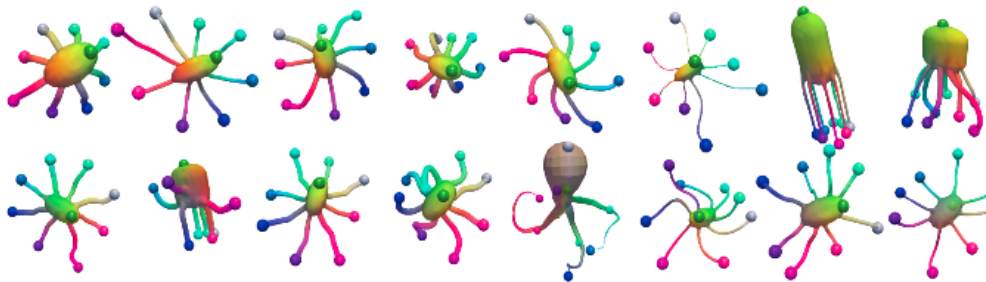


Figure 1: Consistent Collection Matching. Results of the proposed one-stage procedure for finding consistent correspondences between shapes in a collection showing strong variability and non-rigid deformations.

Abstract

Establishing point correspondences between shapes is extremely challenging as it involves both finding sets of semantically persistent feature points, as well as their combinatorial matching. We focus on the latter and consider the Quadratic Assignment Matching (QAM) model. We suggest a novel convex relaxation for this NP-hard problem that builds upon a rank-one reformulation of the problem in a higher dimension, followed by relaxation into a semidefinite program (SDP). Our method is shown to be a certain hybrid of the popular spectral and doubly-stochastic relaxations of QAM and in particular we prove that it is tighter than both.

Experimental evaluation shows that the proposed relaxation is extremely tight: in the majority of our experiments it achieved the certified global optimum solution for the problem, while other relaxations tend to produce sub-optimal solutions. This, however, comes at the price of solving an SDP in a higher dimension.

Our approach is further generalized to the problem of Consistent Collection Matching (CCM) where we solve

$$\begin{aligned} \max_Y \quad & \text{tr}(WY) & (7a) \\ \text{s.t.} \quad & Y \succeq [X][X]^T & (7b) \\ & X \in \text{conv} \Pi_n^k & (7c) \\ & \text{tr}Y = k & (7d) \\ & Y \succeq 0 & (7e) \\ & \sum_{qrst} Y_{qrst} = k^2 & (7f) \\ & Y_{qrst} \leq \begin{cases} 0, & \text{if } q = s, r \neq t \\ 0, & \text{if } r = t, q \neq s \\ \min\{X_{qr}, X_{st}\}, & \text{otherwise} \end{cases} & (7g) \end{aligned}$$

$$\begin{aligned} \max_{X,Y} \quad & \sum_{i,j} \text{tr}(W^{ij}Y^{ij}) & (10a) \\ \text{s.t.} \quad & (X^{ij}, Y^{ij}) \in \mathcal{C}^k \quad \forall i < j & (10b) \\ & X^{ii} \in \mathcal{D} \cap \text{conv} \Pi_n^k \quad \forall i & (10c) \\ & X \succeq 0 & (10d) \end{aligned}$$

Approximate Methods

Consistent Partial Matching of Shape Collections via Sparse Modeling

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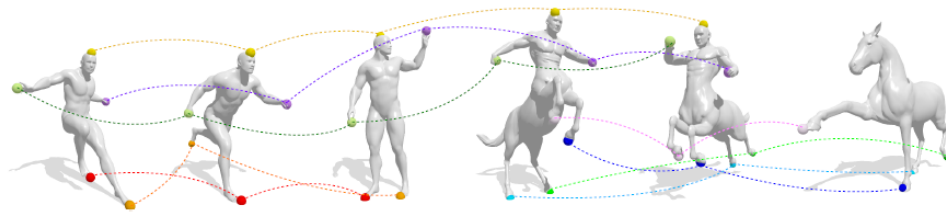


Figure 1: A partial multi-way correspondence obtained with our approach on a heterogeneous collection of shapes. Our method does not require initial pairwise maps as input, as it actively seeks a reliable correspondence in the space of joint, cycle-consistent matches. Partially-similar as well as outlier shapes are automatically filtered out for by adopting a sparse model for the joint correspondence. A subset of all matches is shown.

Abstract

Recent efforts in the area of joint object matching approach the problem by taking as input a set of pairwise maps which are then jointly optimized across the whole collection so that certain accuracy is satisfied. One natural requirement is cycle-consistency – namely the fact that map composition yields the same result regardless of the path taken in the shape collection. In this paper, we introduce a method to obtain consistent matches without requiring initial pairwise solutions to be given as input. We define a joint measure of metric distortion directly over the space of cycle-consistent maps; in particular, for similar and extra-class shapes, we formulate the problem as a series of quadratic programming constraints, making our technique a natural candidate for analyzing collections with outliers. The particular form of the problem allows us to leverage results and tools from the theory. This enables a highly efficient optimization procedure which assures accurate solutions in a matter of minutes in collections with hundreds of shapes.

Categories and Subject Descriptors (according to ACM CCS): I.3.5 [Computer Graphics] Visualization and Object Modeling—Shape Analysis

1. Introduction

Finding matches among multiple objects is a research topic

this end, a natural and widely accepted criterion is cycle-consistency [ZKP10], namely that composition of maps

Sequence of quadratic programs; based on metric distortion and WKS descriptor match

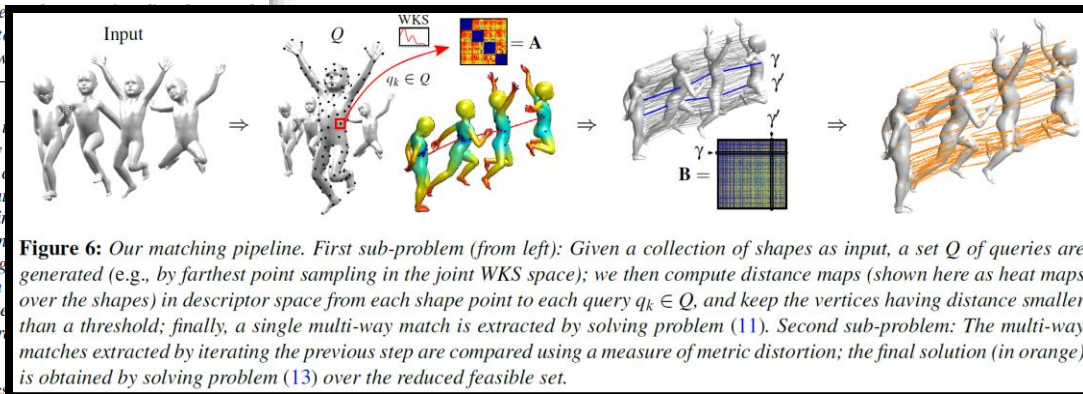
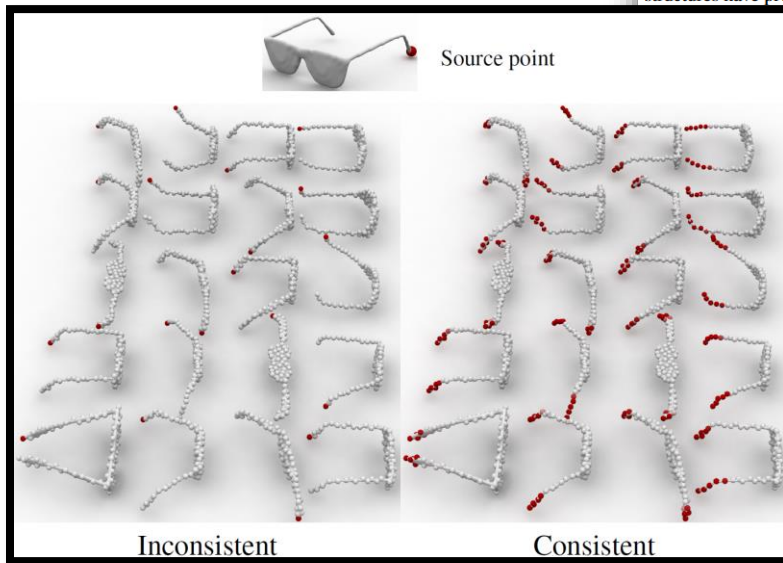


Figure 6: Our matching pipeline. First sub-problem (from left): Given a collection of shapes as input, a set Q of queries are generated (e.g., by farthest point sampling in the joint WKS space); we then compute distance maps (shown here as heat maps over the shapes) in descriptor space from each shape point to each query $q_k \in Q$, and keep the vertices having distance smaller than a threshold; finally, a single multi-way match is extracted by solving problem (11). Second sub-problem: The multi-way matches extracted by iterating the previous step are compared using a measure of metric distortion; the final solution (in orange) is obtained by solving problem (13) over the reduced feasible set.

Approximate Methods

Multiplicative updates
for nonconvex
nonnegative matrix
factorization



$$\min_A \text{KL}(G | AA^T)$$

Entropic Metric Alignment for Correspondence Problems

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Abstract

Many shape and image processing tools rely on computation of correspondences between geometric domains. Efficient methods that stably extract “soft” matches in the presence of diverse geometric structures have proven to be valuable for shape retrieval and transfer of information. With these applications in mind, we propose an algorithm for probabilistic correspondence that optimizes an entropically-regularized Gromov-Wasserstein (GW) objective. Our developments in numerical optimal transportation, which are compact, provably convergent, and applicable to a wide range of domains, can be expressed as a metric measure matrix. We present extensive experiments illustrating the convergence of our algorithm to a variety of graphics tasks. We expand entropic GW correspondence to a framework for matching problems, incorporating partial distance matrices, shape exploration, symmetry detection, and more than two domains. These applications expand entropic GW correspondence to major shape analysis tasks, including robustness to distortion and noise.

Gromov-Wasserstein, matching, entropy

Computing methodologies → Shape analysis;

Con

One of the geometry processing toolbox is a tool for correspondence, the problem of finding which points on one shape correspond to points on a source. Many variations of this problem have been considered in the graphics literature, e.g. point-to-point matching with some sparse correspondences provided by the user. Regardless, the basic task of geometric correspondence facilitates the transfer of properties and edits from one shape to another.

The primary factor that distinguishes correspondence algorithms is the choice of objective functions. Different choices of objective functions express contrasting notions of which correspondences are “desirable.” Classical theorems from differential geometry and most modern algorithms consider *local* distortion, producing maps that take tangent planes to tangent planes with as little stretch as possible; slightly larger neighborhoods might be taken into account by e.g.

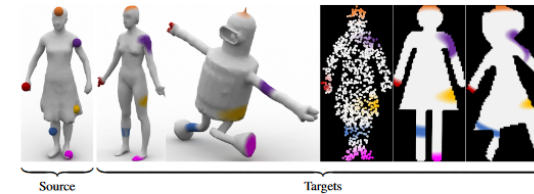


Figure 1: Entropic GW can find correspondences between a source surface (left) and a surface with similar structure, a surface with shared semantic structure, a noisy 3D point cloud, an icon, and a hand drawing. Each fuzzy map was computed using the same code.

These algorithms suffer from having to patch together local elastic terms into a single global map.

In this paper, we propose a new correspondence algorithm that minimizes distortion of long- and short-range distances alike. We study an entropically-regularized version of the *Gromov-Wasserstein* (GW) mapping objective function from [Mémoli 2011] measuring the distortion of geodesic distances. The optimizer is a probabilistic matching expressed as a “fuzzy” correspondence matrix in the style of [Kim et al. 2012; Solomon et al. 2012]; we control sharpness of the correspondence via the weight of an entropic regularizer.

Although [Mémoli 2011] and subsequent work identified the possibility of using GW distances for geometric correspondence, computational challenges hampered their practical application. To overcome these challenges, we build upon recent methods for regularized optimal transportation introduced in [Benamou et al. 2015; Solomon et al. 2015]. While optimal transportation is a fundamentally different optimization problem from regularized GW computation (linear versus quadratic matching), the core of our method relies upon solving a sequence of regularized optimal transport problems.

Our remarkably compact algorithm (see Algorithm 1) exhibits global convergence, i.e., it *provably* reaches a local minimum of the regularized GW objective function regardless of the initial guess. Our algorithm can be applied to any domain expressible as a metric measure space (see §2). Concretely, only distance matrices are required as input, and hence the method can be applied to many classes of domains including meshes, point clouds, graphs, and even more

Computer Vision Perspective

Unpaired Image-to-Image Translation using Cycle-Consistent Adversarial Networks

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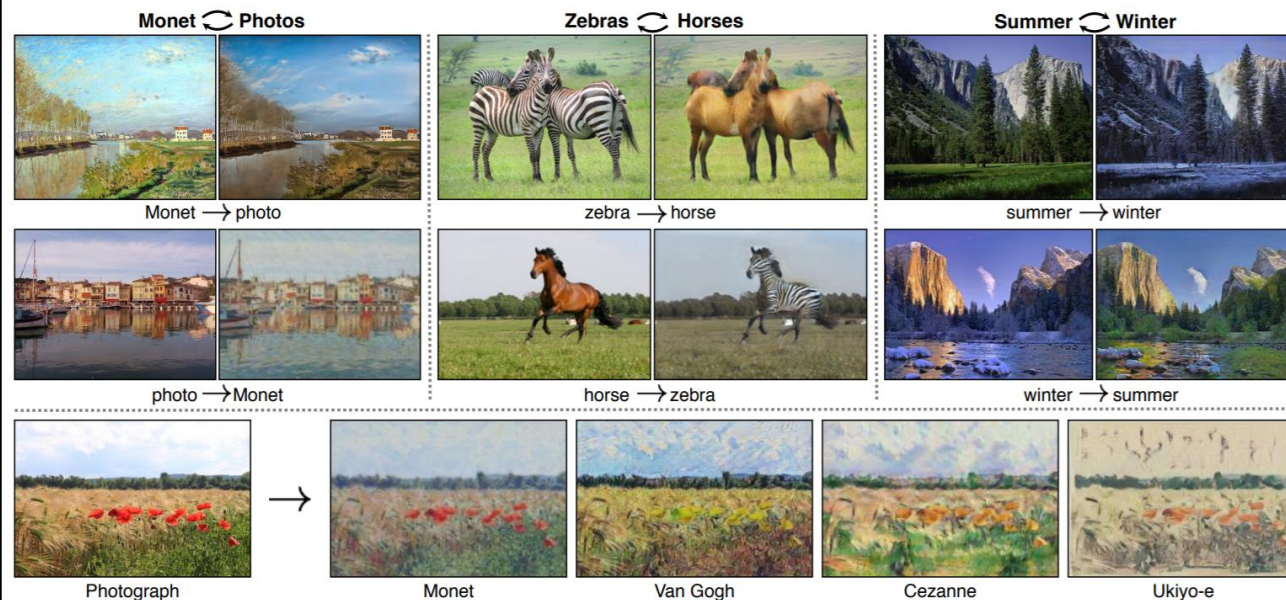


Figure 1: Given any two unordered image collections X and Y , our algorithm learns to automatically “translate” an image from one into the other and vice versa: (left) Monet paintings and landscape photos from Flickr; (center) zebras and horses from ImageNet; (right) summer and winter Yosemite photos from Flickr. Example application (bottom): using a collection of paintings of famous artists, our method learns to render natural photographs into the respective styles.

Abstract

Image-to-image translation is a class of vision and graphics problems where the goal is to learn the mapping

1. Introduction

What did Claude Monet see as he placed

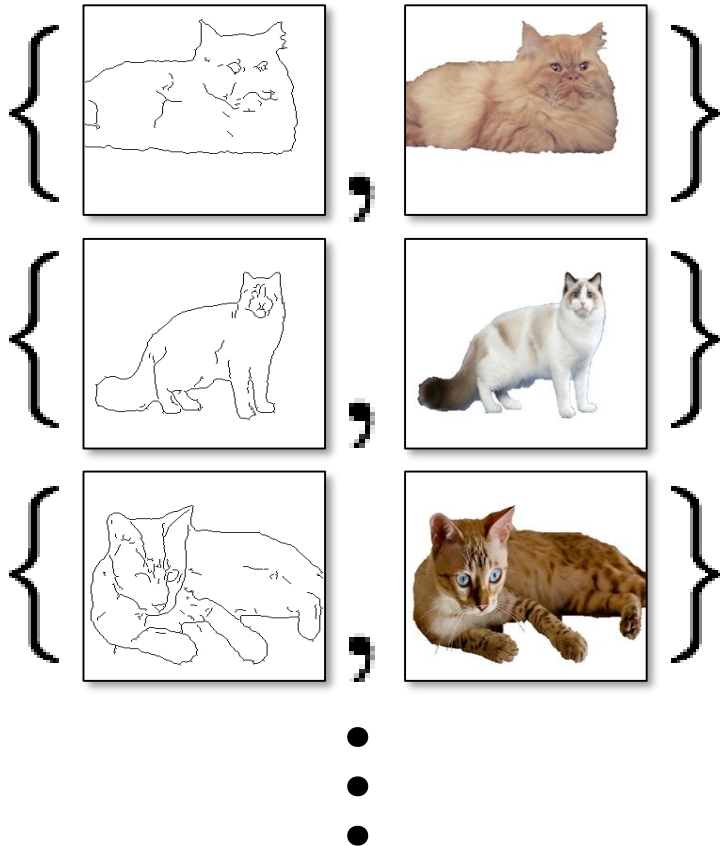
Slides courtesy the authors
<https://junyanz.github.io/CycleGAN/>

Paired vs. Unpaired Problems

Paired

x_i

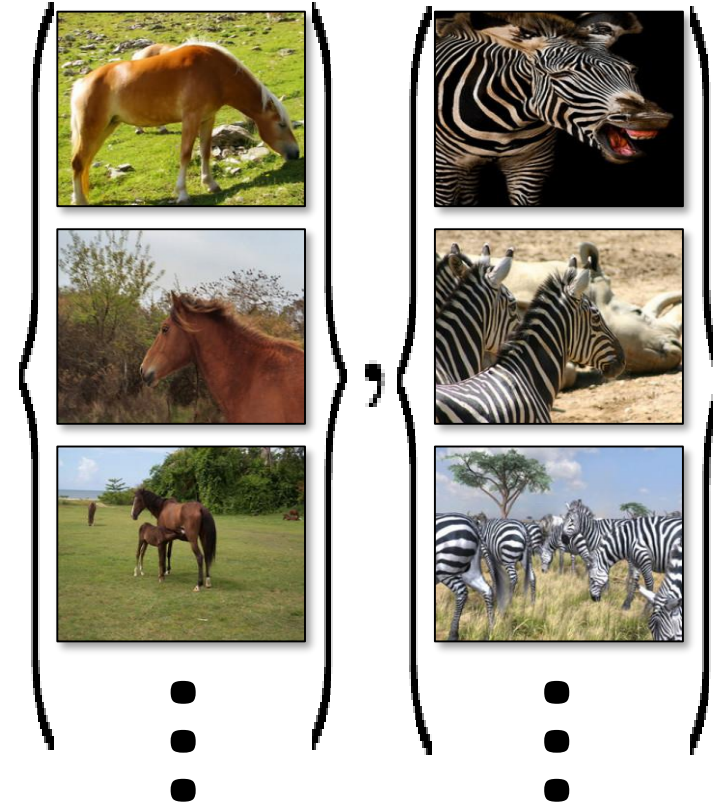
y_i



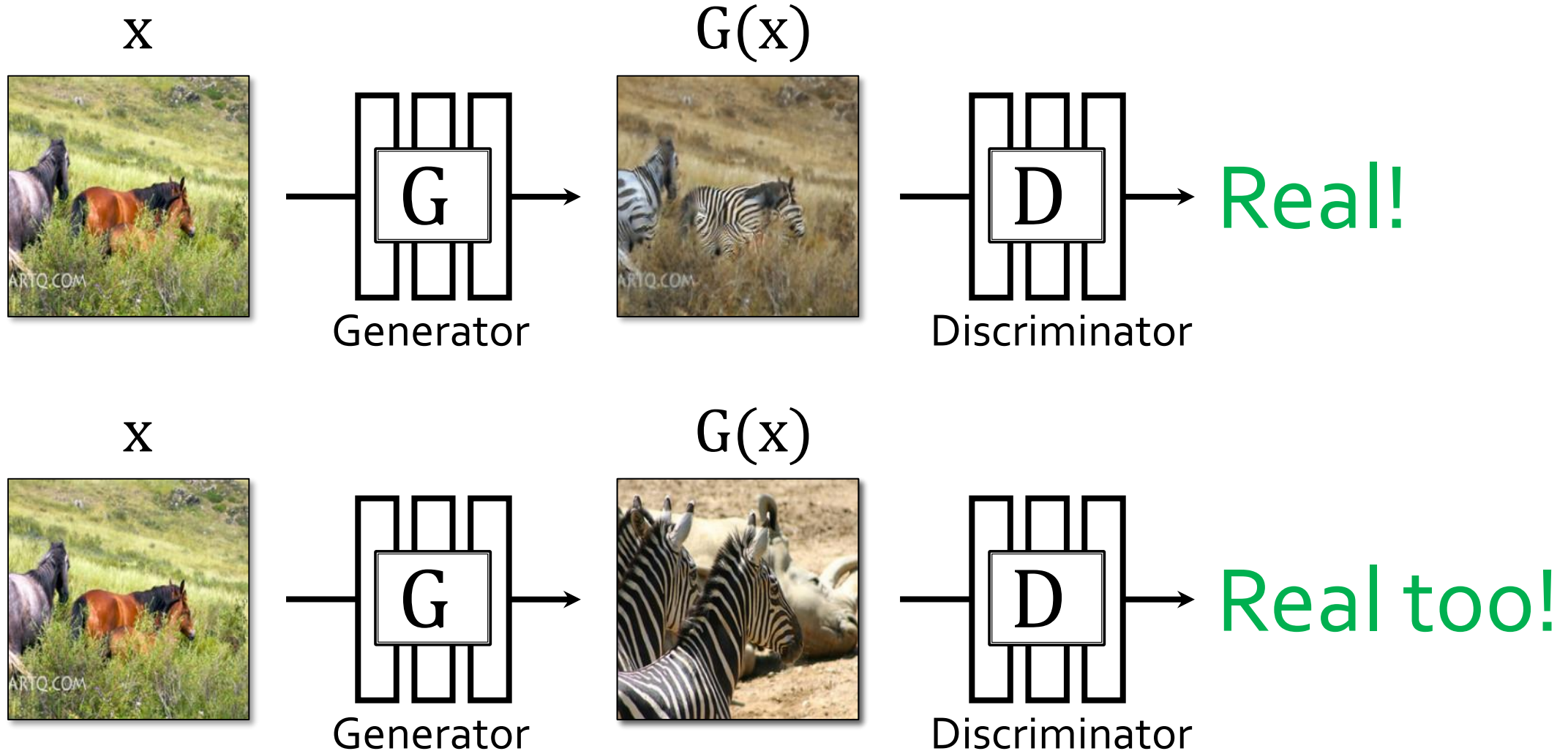
Unpaired

X

Y



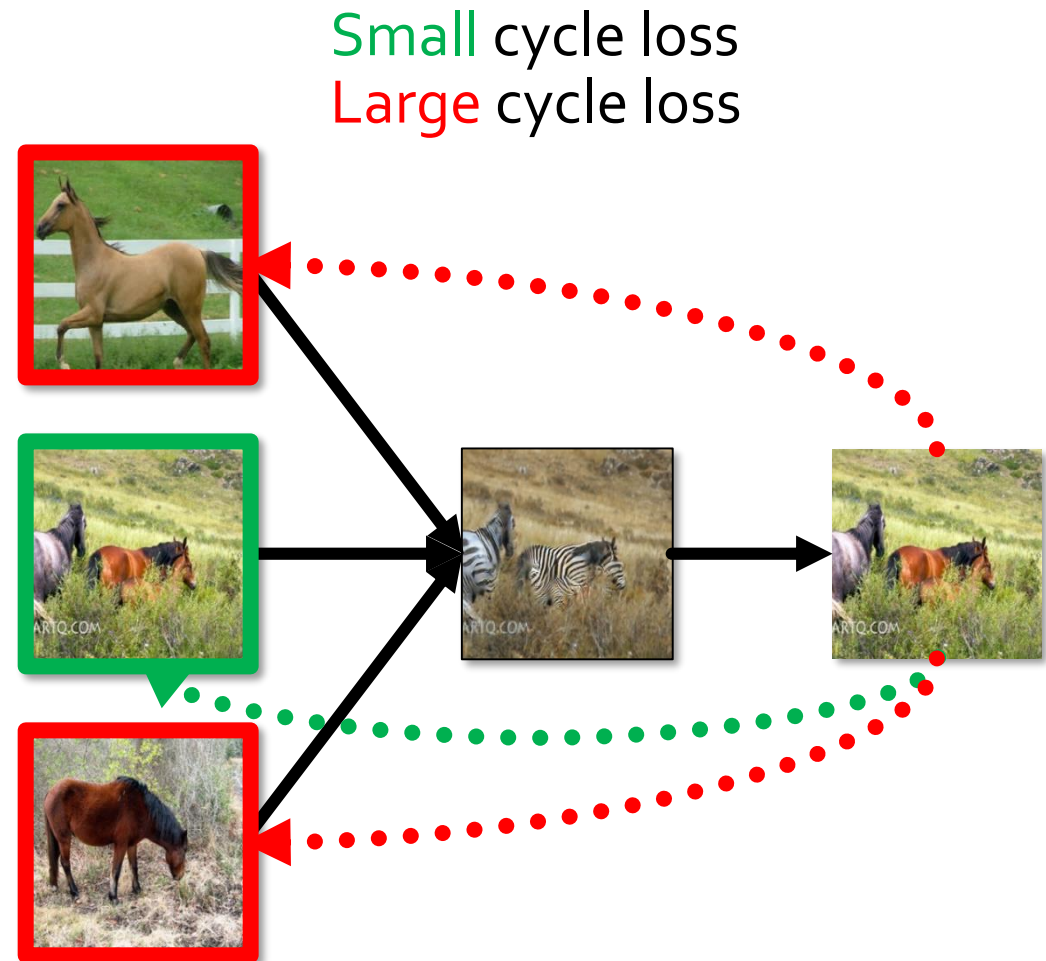
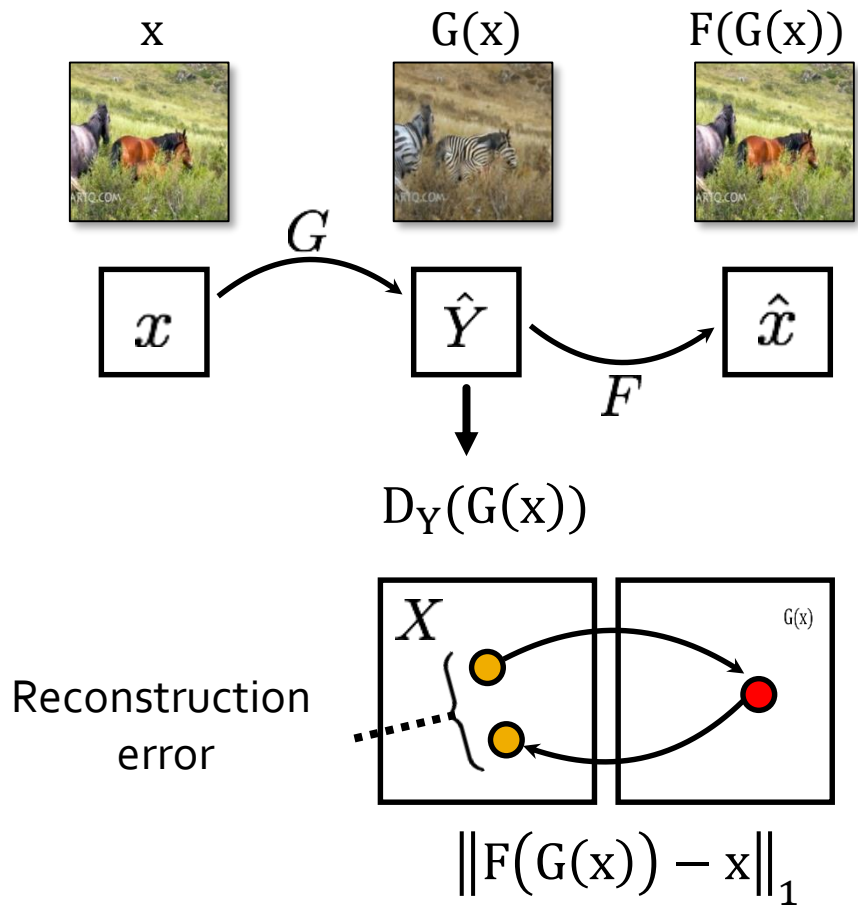
Adversarial Networks: Problem



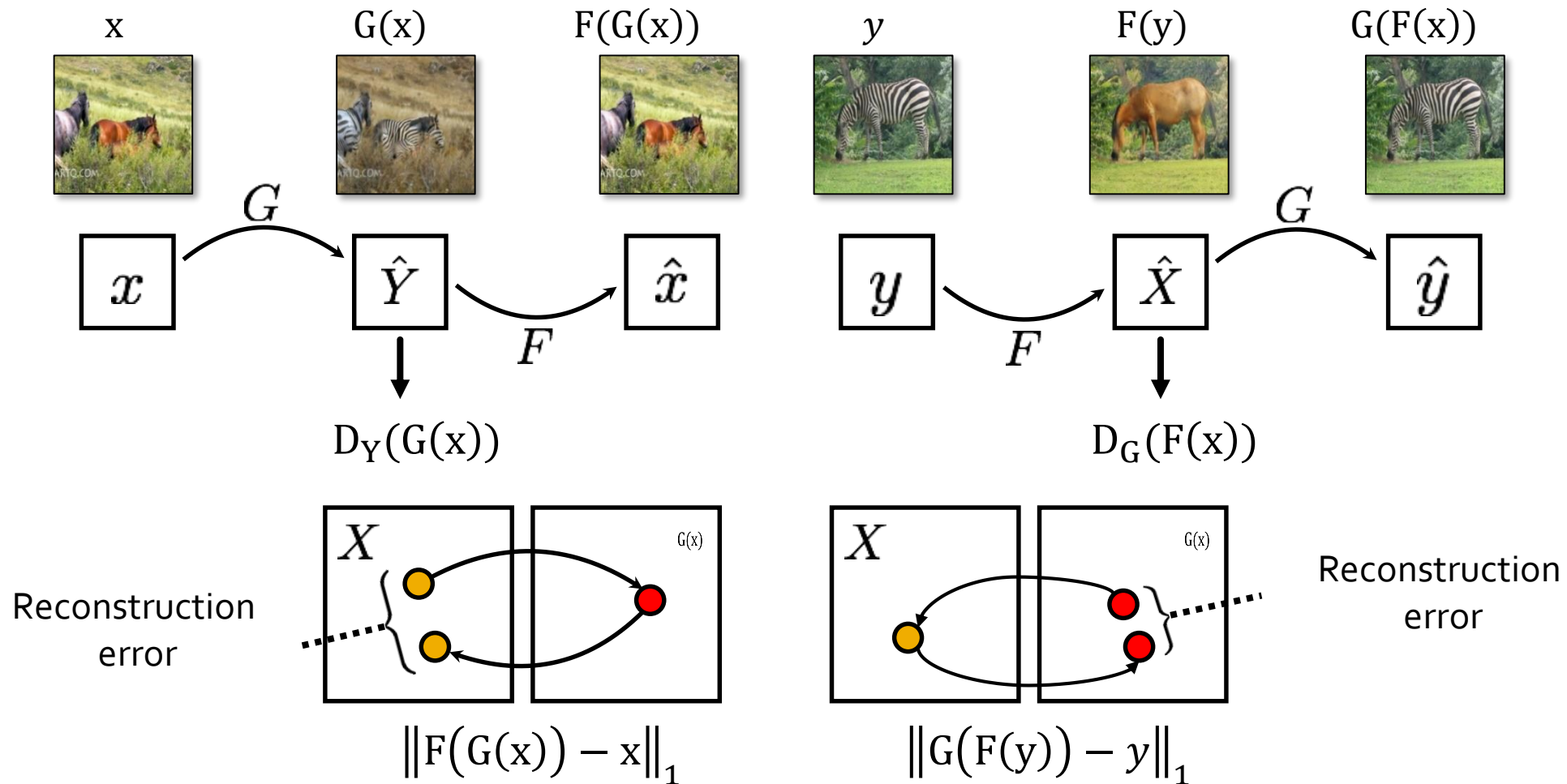


Mode collapse

Cycle-Consistent Adversarial Networks



Cycle Consistency Loss







Input

Monet

Van Gogh

Cezanne

Ukiyo-e

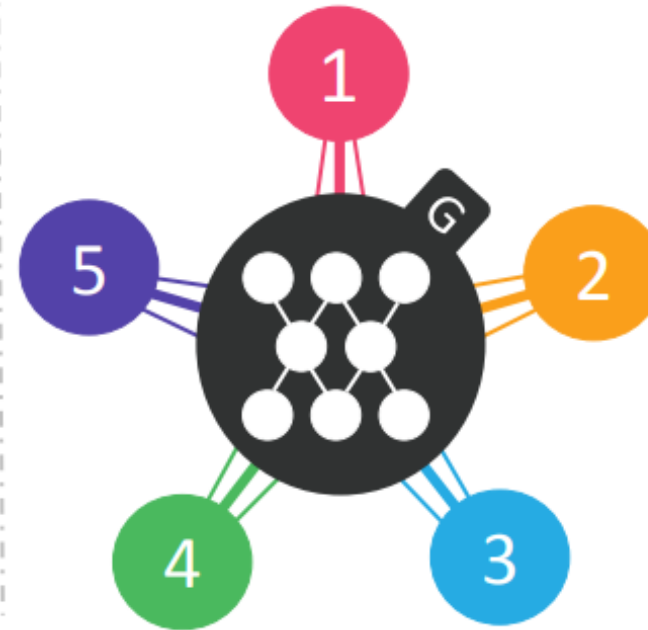


More Than Two Domains?

(a) Cross-domain models



(b) StarGAN



Consistent Correspondence

Justin Solomon

6.8410: Shape Analysis

Spring 2023



Extra: Angular Synchronization

Justin Solomon

6.8410: Shape Analysis

Spring 2023

