

High-Productivity Stream Programming for High-Performance Systems

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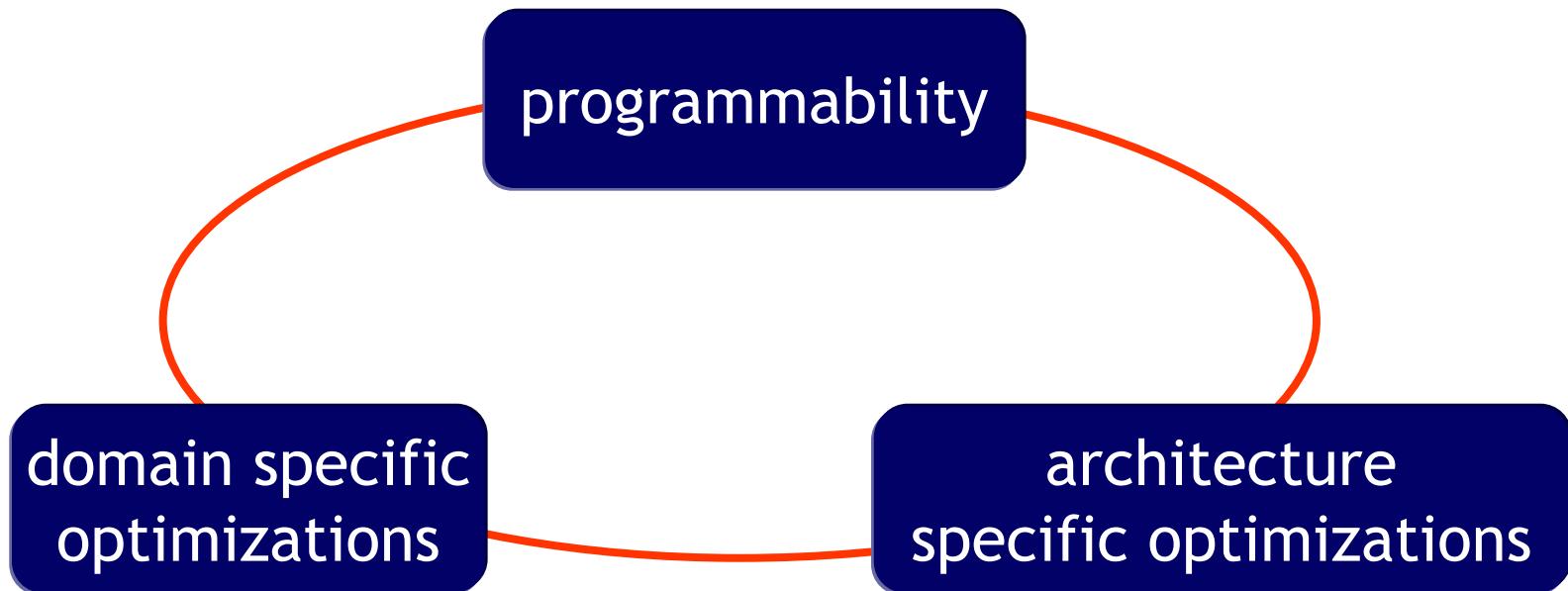
Massachusetts Institute of Technology



<http://cag.lcs.mit.edu/streamit>

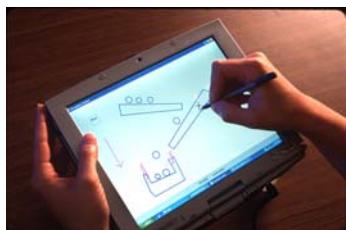
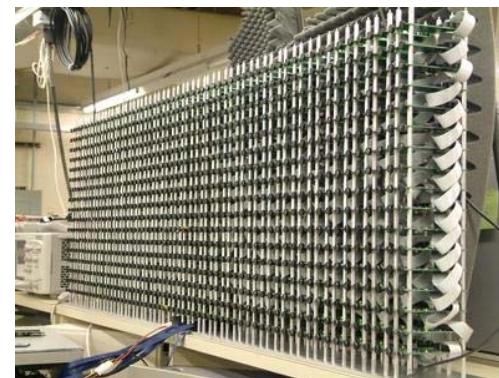
The StreamIt Vision

- Boost productivity, enable faster development and rapid prototyping



- Simple and effective optimizations for streams
- Targeting tiled architectures, clusters of workstations, DSPs, and traditional uniprocessors

Why an Emphasis on Streaming?

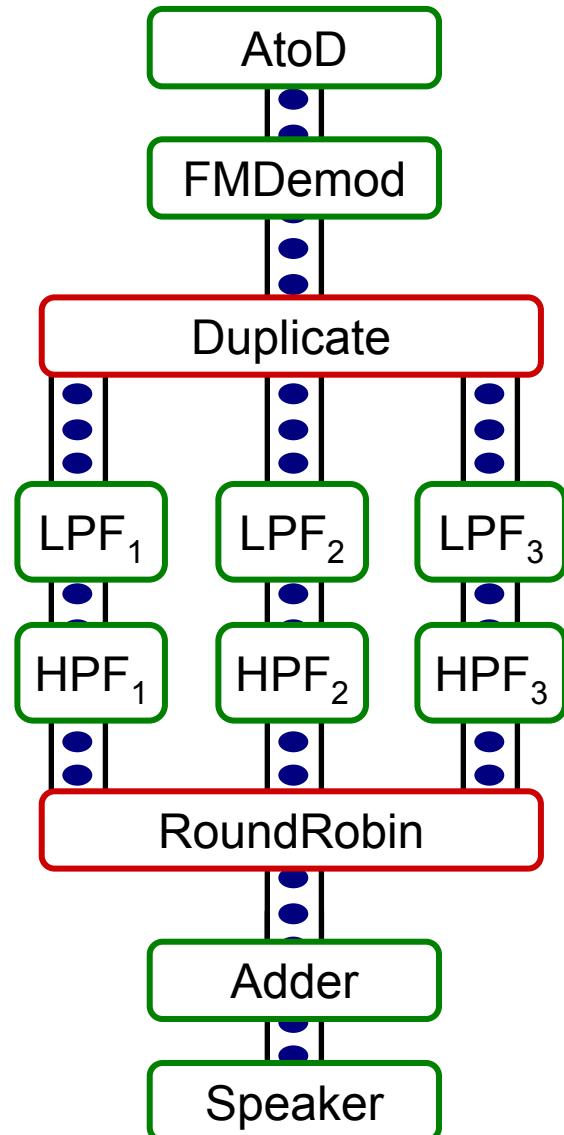


Streaming in other Domains as well

- Cryptography
- Databases
- Face recognition
- Network processing and security
- Scientific codes
- ...
- Attractive programming model because of a simple mapping from specification to implementation

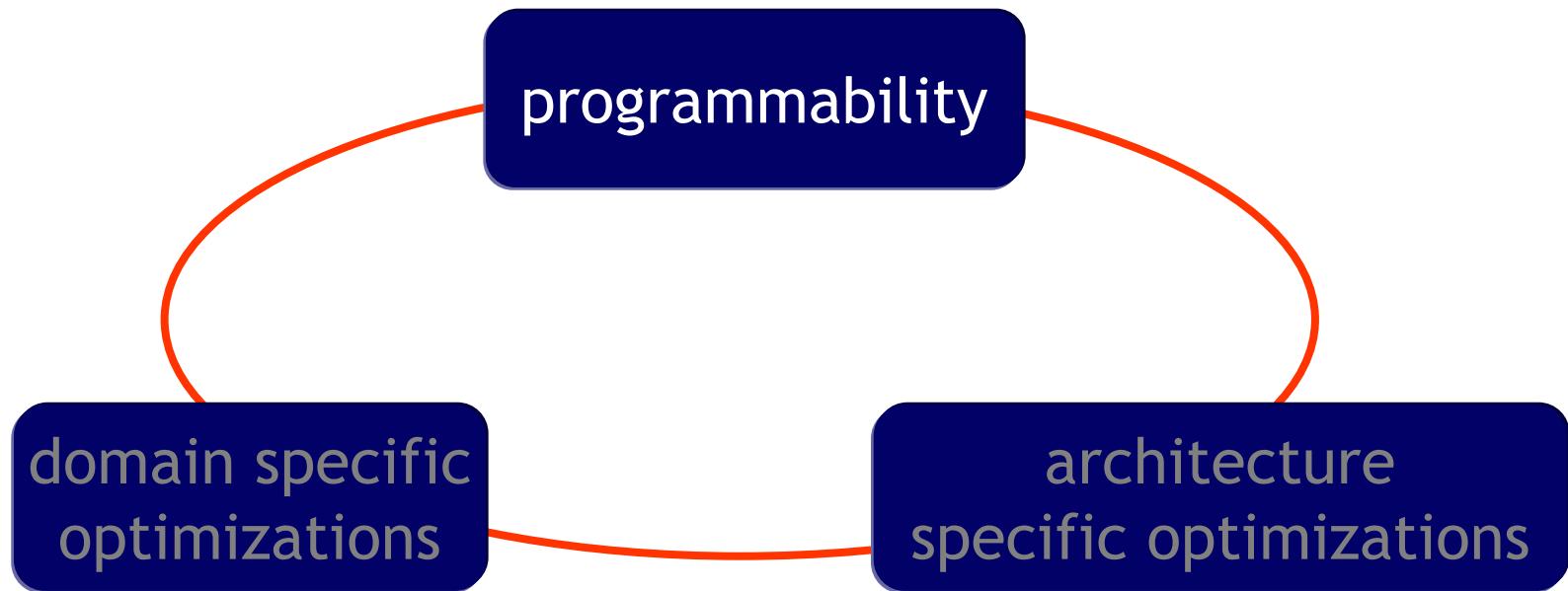
Properties of Stream Programs

- Mostly regular and repeating computation
- Parallel, independent computation with explicit communication
- Amenable to aggressive compiler optimizations
[ASPLOS '02, PLDI '03, LCTES'03, LCTES '05]



The StreamIt Vision

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- Simple and effective optimizations for streams
- Targeting tiled architectures, clusters of workstations, DSPs, and traditional uniprocessors

Programming in StreamIt

```
void->void pipeline FMRadio(int N, float freq1, float freq2) {
```

```
    add AtoD();
```

- Natural correspondence
between text and
application graph

```
    add FMDemod();
```

```
    add splitjoin {
        split duplicate;
        for (int i=0; i<N; i++) {
            add pipeline {
```

```
                add LowPassFilter(freq1 + i*(freq2-freq1)/N);
```

```
                add HighPassFilter(freq2 + i*(freq2-freq1)/N);
```

```
            }
        }
```

```
        join roundrobin();
```

```
}
```

```
    add Adder();
```

```
    add Speaker();
```

Programming in StreamIt

```
void->void pipeline FMRadio(int N, float freq1, float freq2) {
```

```
    add AtoD();
```

- Streams are easily composed

```
    add FMDemod();
```

```
    add splitjoin {
        split duplicate;
        for (int i=0; i<N; i++) {
            add pipeline {
```

```
                add LowPassFilter(freq1 + i*(freq2-freq1)/N);
```

```
                add HighPassFilter(freq2 + i*(freq2-freq1)/N);
```

```
            }
        }
```

```
        join roundrobin();
```

```
}
```

```
    add Adder();
```

```
    add Speaker();
```

Programming in StreamIt

```
void->void pipeline FMRadio(int N, float freq1, float freq2) {
```

```
    add AtoD();
```

- Streams are parameterized, and malleable

```
    add FMDemod();
```

```
    add splitjoin {
        split duplicate;
        for (int i=0; i<N; i++) {
            add pipeline {
```

```
                add LowPassFilter(freq1 + i*(freq2-freq1)/N);
```

```
                add HighPassFilter(freq2 + i*(freq2-freq1)/N);
```

```
            }
        }
```

```
        join roundrobin();
```

```
}
```

```
    add Adder();
```

```
    add Speaker();
```

Programming in StreamIt

```
void->void pipeline FMRadio(int N, float freq1, float freq2) {
```

```
    add AtoD();
```

- Application is
architecture independent
(i.e., portable)

```
    add FMDemod();
```

```
    add splitjoin {
        split duplicate;
        for (int i=0; i<N; i++) {
            add pipeline {
```

```
                add LowPassFilter(freq1 + i*(freq2-freq1)/N);
```

```
                add HighPassFilter(freq2 + i*(freq2-freq1)/N);
```

```
            }
        }
```

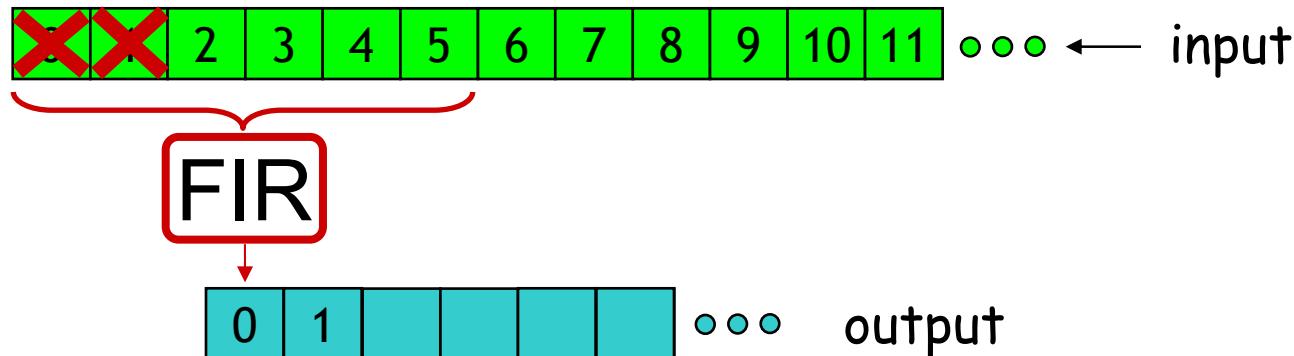
```
        join roundrobin();
```

```
}
```

```
    add Adder();
```

```
    add Speaker();
```

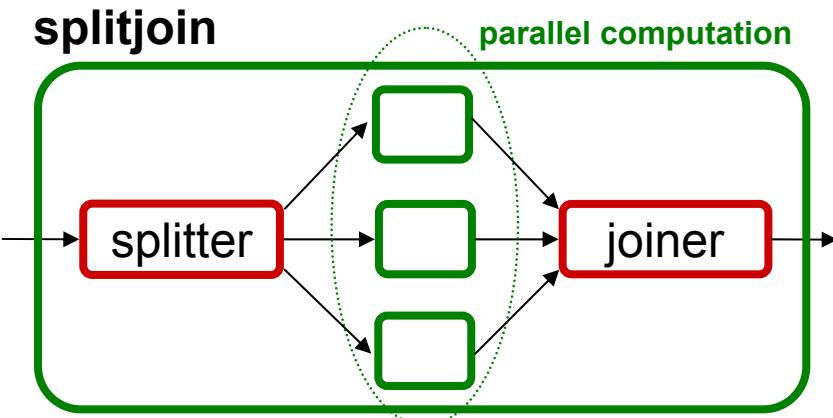
Filters as Computational Elements



```
float→float filter FIR (int N) {
    work push 1 pop 1 peek N {
        float result = 0;
        for (int i = 0; i < N; i++) {
            result += weights[i] * peek(i);
        }
        push(result);
        pop();
    }
}
```

Benefits of StreamIt

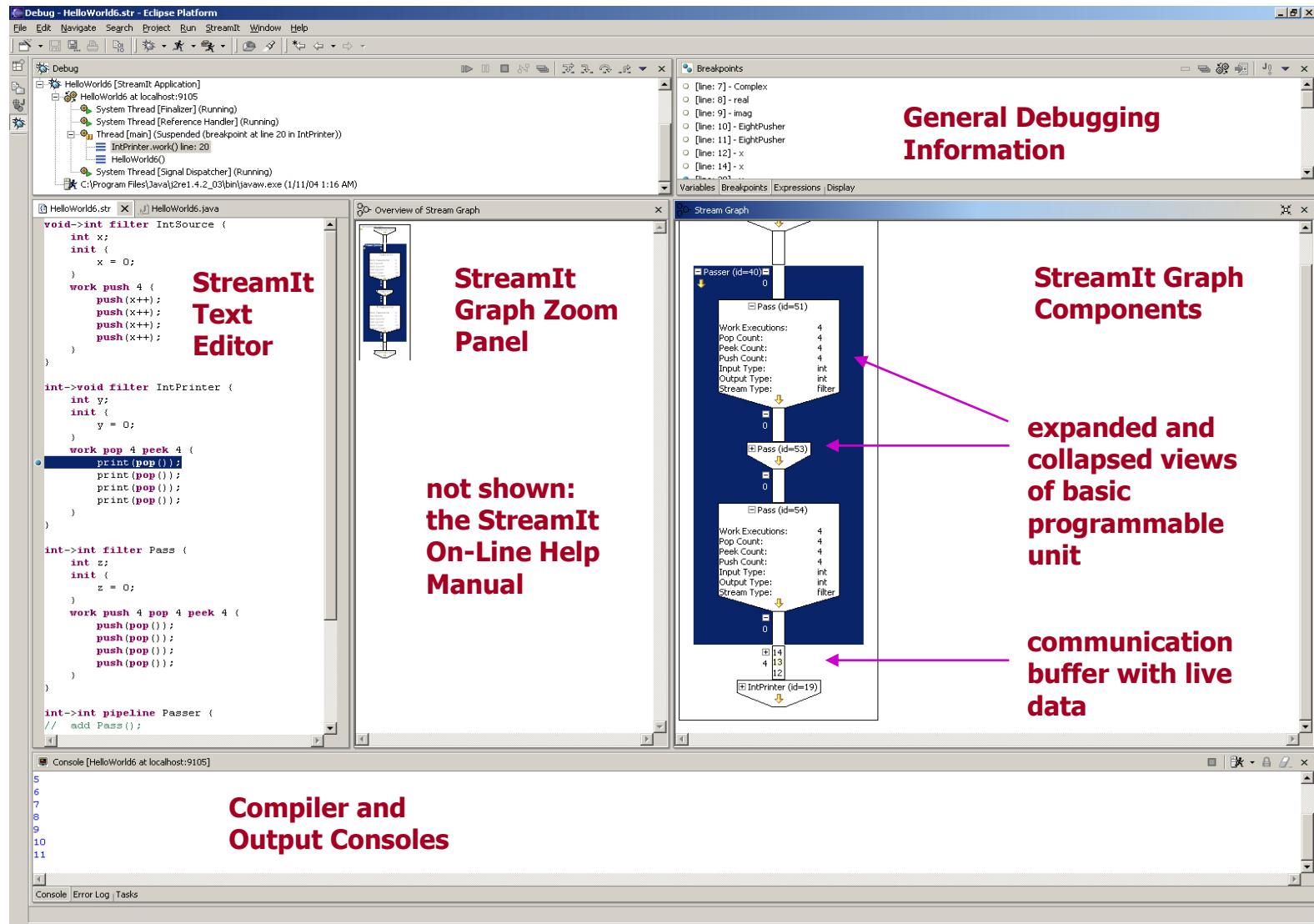
- Communication is exposed and pipeline parallelism is more readily discovered
- Flow of data provides a frame of reference for reasoning about “time” [PPoPP ’05]
 - Powerful advantage when debugging parallel programs



versus

- Multiple threads with independent program counters
- Non-deterministic execution

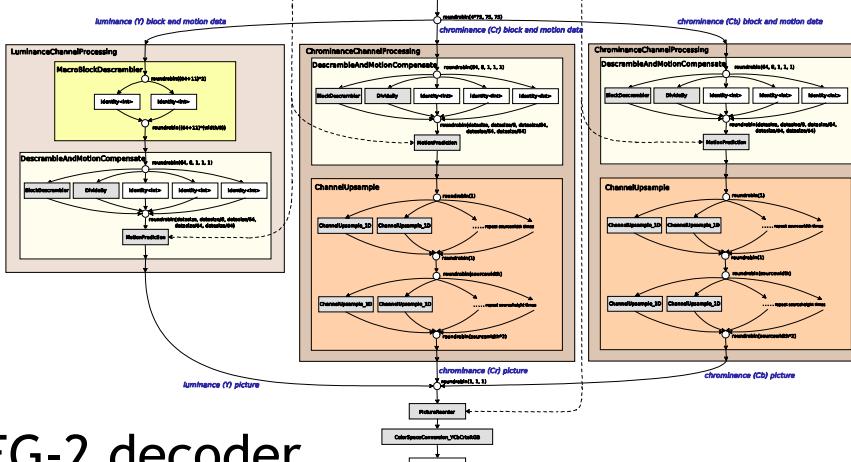
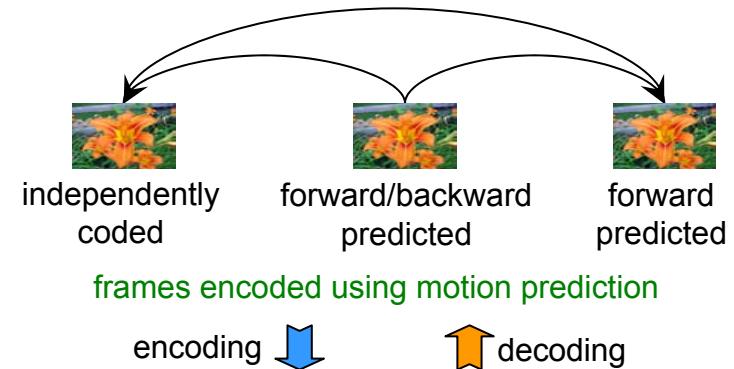
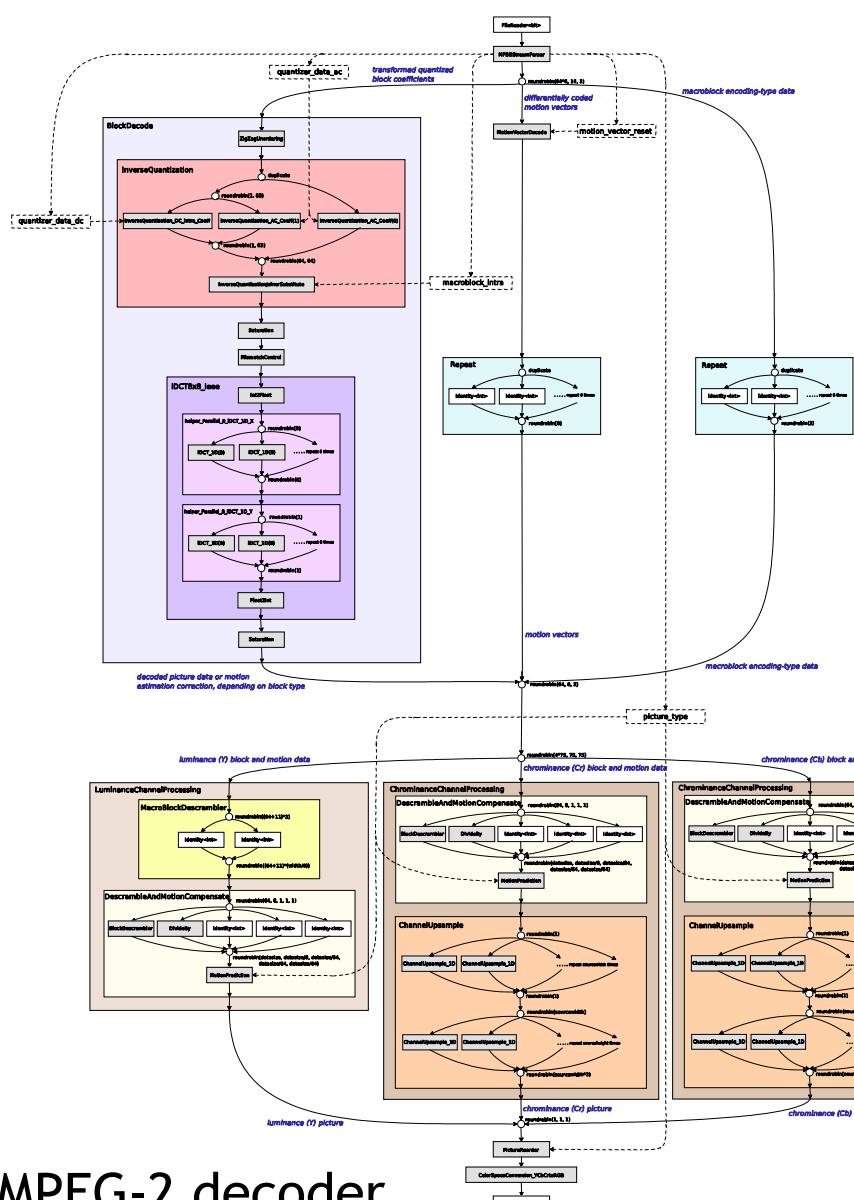
StreamIt Development Environment



StreamIt Applications

- Software radio
- Frequency hopping radio
- Acoustic beam former
- Vocoder
- GMTI (ground moving target indicator)
- DES and Serpent blocked ciphers
- Sorting
- FFTs and DCTs
- JPEG
- ...

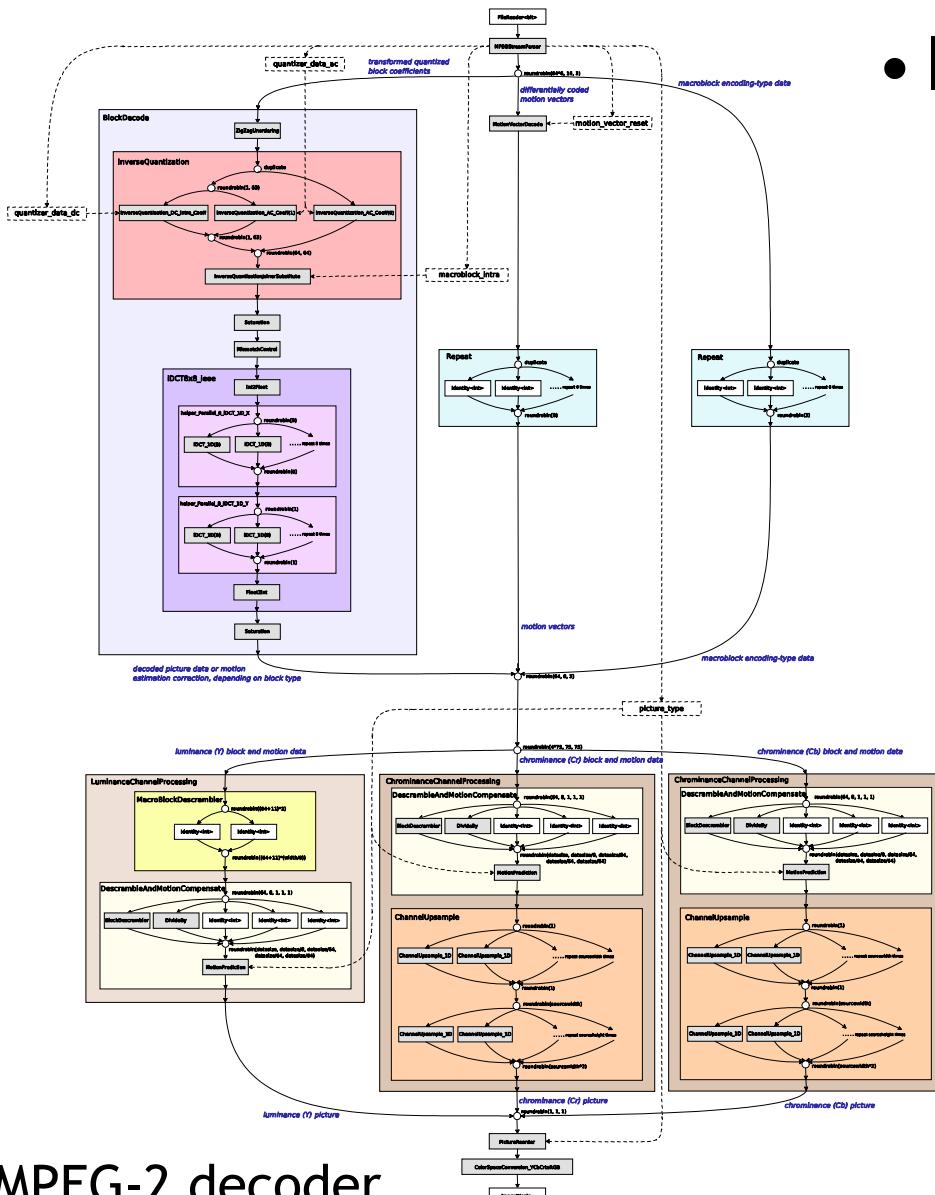
MPEG: Motion Video Codec



-1	1	1	1	0	0	0	0
0	1	0	0	0	0	0	0
1	0	0	0	0	0	0	0
-2	0	0	0	0	0	0	0
0	3	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0

DCT and quantization of 8x8 image block

MPEG: Motion Video Codec



- Implementation statistics
 - 4921 lines of code
 - 48 static streams
 - Compile to ~2150 filters
 - 352x240 resolution
- Reference C implementation has 9832 lines of code
 - Supports interlacing and multi-layer streams
- 8 weeks of development
- 1 programmer with no prior MPEG-2 experience

Excerpt from StreamIt Implementation

Specification in Section 7.4.1: $F''[0][0] = \text{intra_dc_mult} \times QF[0][0]$

Table 7-4 - Relation between intra_dc_precision and intra_dc_mult

intra_dc_precision	bits_of_precision	intra_dc_mult
0	8	8
1	9	4
2	10	2
3	11	1

```
int->int filter InverseQuantization() {
    int[4] intra_dc_mult = {8, 4, 2, 1};
    int intra_dc_precision;

    work pop 1 push 1 {
        push(intra_dc_mult[intra_dc_precision] * pop());
    }
}
```

Excerpt from Reference Implementation

Specification in Section 7.4.1: $F''[0][0] = \text{intra_dc_mult} \times \text{QF}[0][0]$

Table 7-4 - Relation between intra_dc_precision and intra_dc_mult

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```

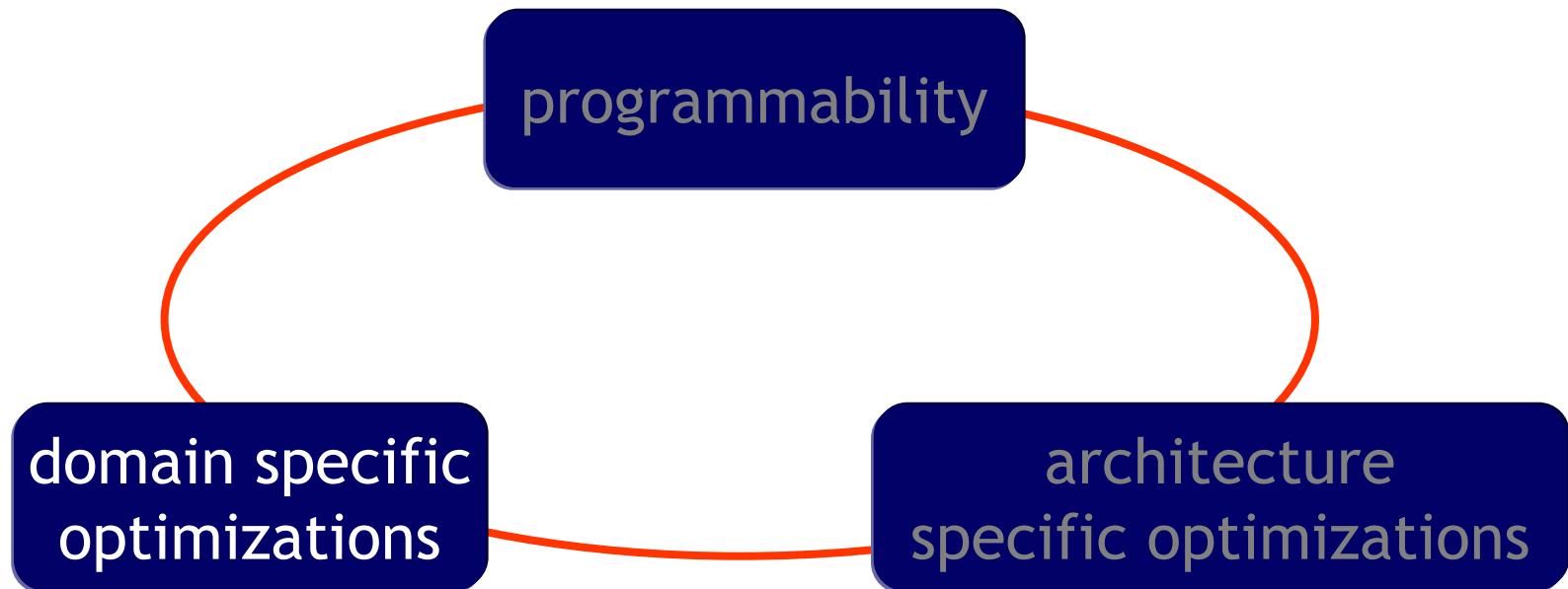
int[4] intra_dc_mult = {8, 4, 2, 1};

for (int m = 0; m < W*H/(16*16); m++)
// six components for chrominance and luminance
for (int comp = 0; comp < 6; comp++)
if (macroblock[m].intra)
    macroblock[m].block[comp][0] *= intra_dc_mult[intra_dc_precision];
// and many lines later
if (cc == 0)
    val = (dc_dct_pred[0] += Get_Luma_DC_dct_diff());
else if (cc == 1)
    val = (dc_dct_pred[1] += Get_Chroma_DC_dct_diff());
else
    val = (dc_dct_pred[2] += Get_Chroma_DC_dct_diff());
if (Fault_Flag) return;
bp[0] = val << (3-intra_dc_precision);

```

The StreamIt Vision

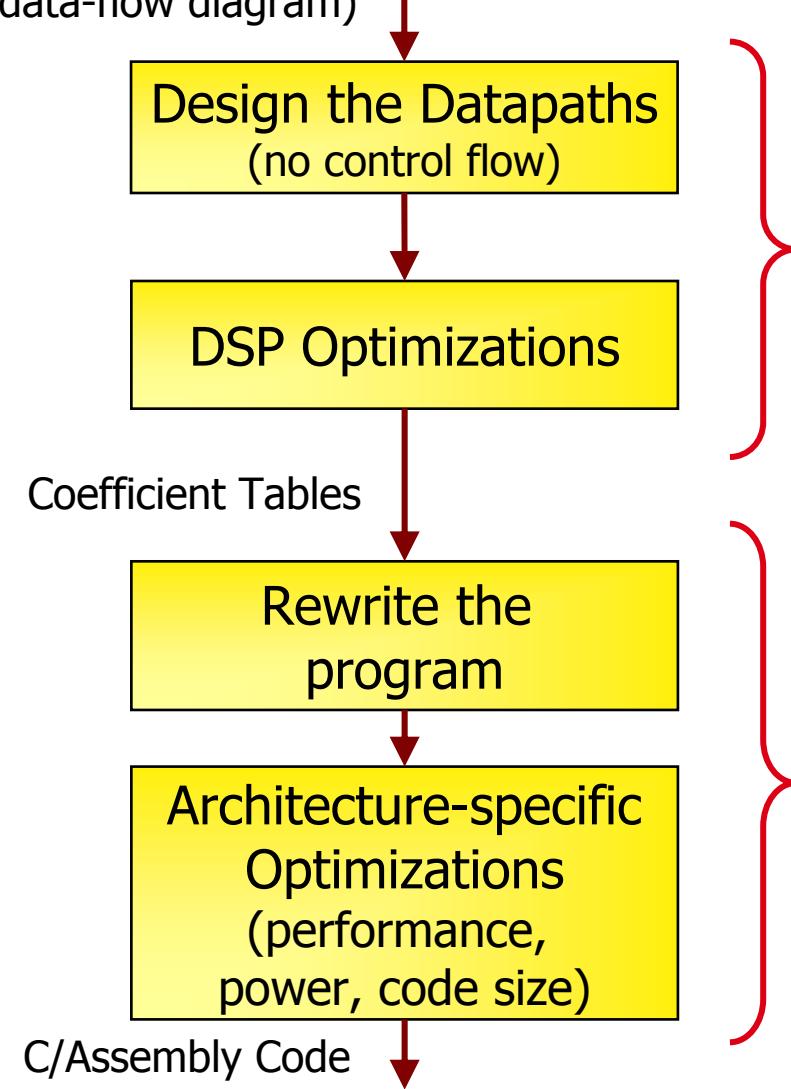
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Conventional DSP Design Flow

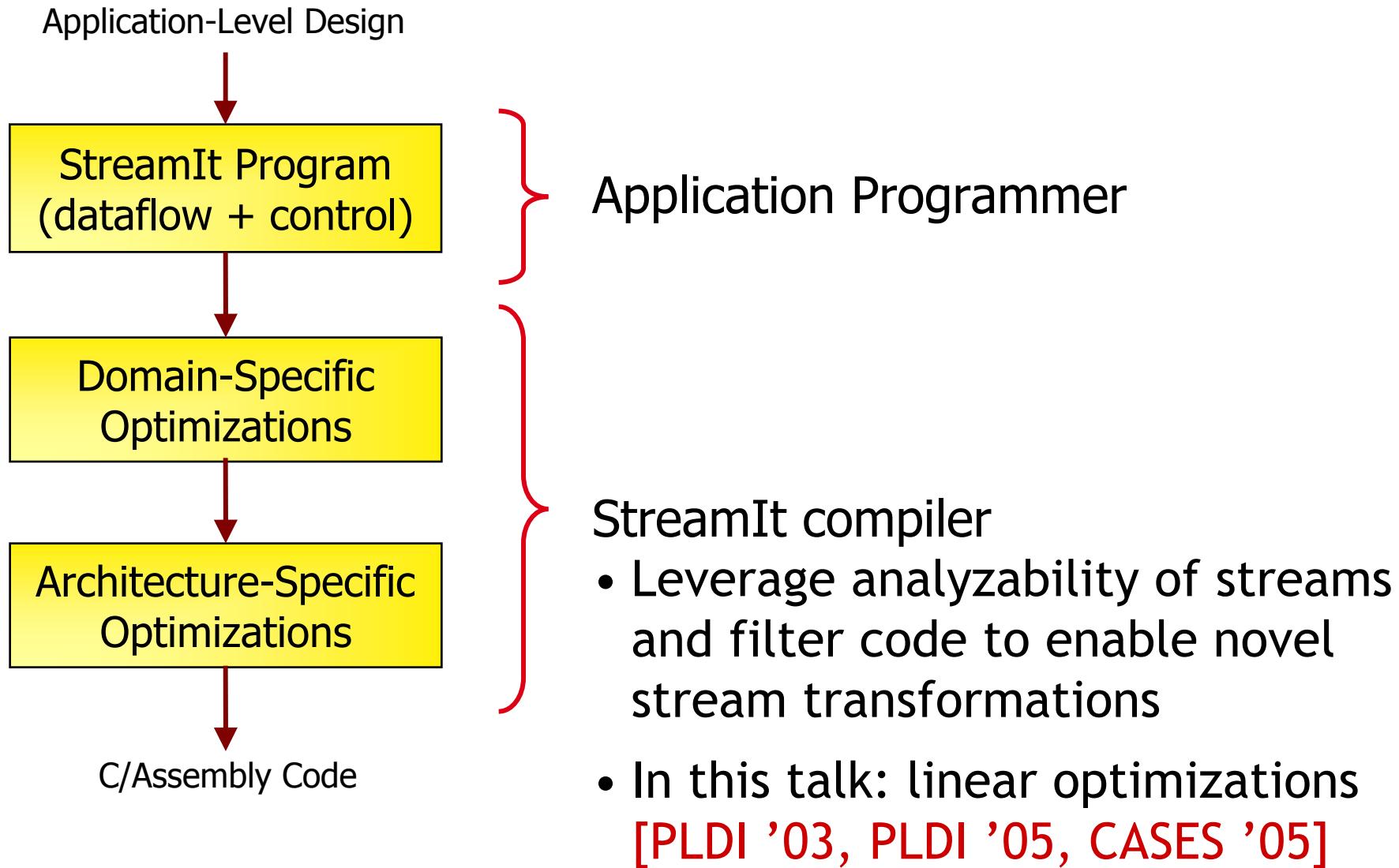
Specification
(data-flow diagram)



Signal Processing Expert
in Matlab

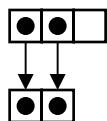
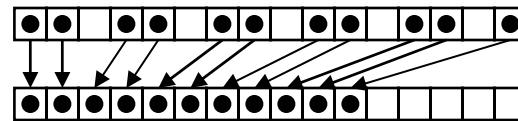
Software Engineer
in C and Assembly

Design Flow with StreamIt



Linear Filter Example

- “Drop every third bit in the bit stream”

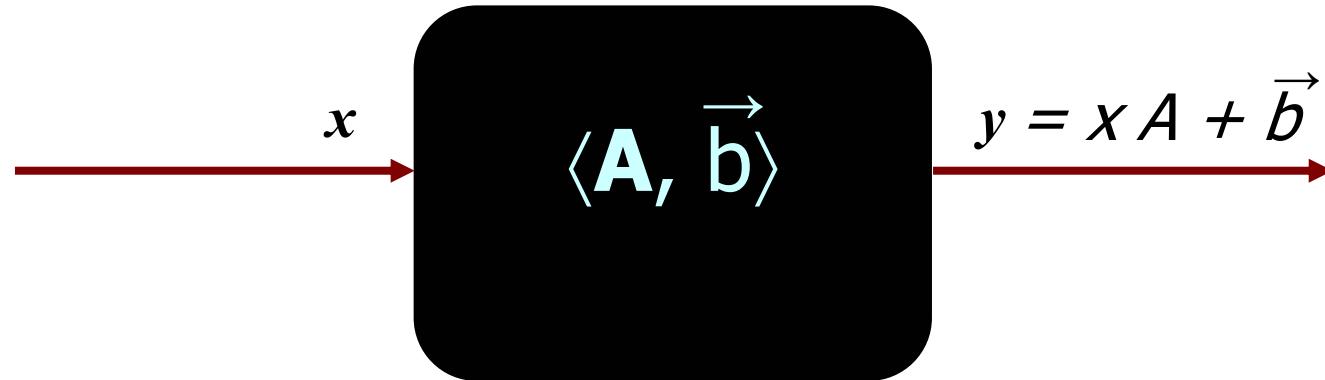


```
bit → bit filter DropThirdBit {
    work push 2 pop 3 {
        push(pop());
        push(pop());
        pop();
    }
}
```

A diagram illustrating a linear filter operation. A 2x3 matrix is multiplied by a vector $\begin{bmatrix} x \\ y \\ z \end{bmatrix}$. The result is a vector $\begin{bmatrix} x \\ y \end{bmatrix}$. The third column of the matrix is highlighted with a green bracket and labeled '3'. The first two columns are labeled '2'.

In General

- A linear filter is a tuple $\langle \mathbf{A}, \vec{\mathbf{b}} \rangle$
 - \mathbf{A} : matrix of coefficients
 - $\vec{\mathbf{b}}$: vector of constants
- Example



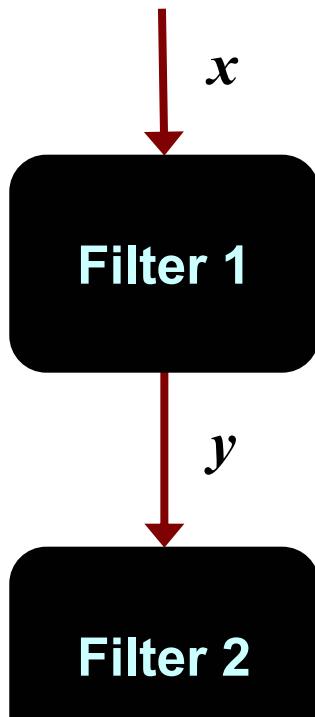
- Linear dataflow analysis resembles constant propagation

Opportunities for Linear Optimizations

- Occur frequently in streaming codes
 - FIR filters
 - Compressors
 - Expanders
 - DFT/DCT
 - Bit permutations in encryption algorithms
 - JPEG and MPEG codecs
 - ...
- Example optimizations
 - Combining adjacent nodes
 - Also, translating to frequency domain when profitable

Combining Linear Filters

$\frac{6 \text{ mults}}{\text{output}}$



$\frac{1 \text{ mult}}{\text{output}}$



$$A = [4 \quad 5 \quad 6]$$

$$y = x A$$

$$z = x \underbrace{A \quad B}_{C}$$

Combined Filter

$$z = x C$$

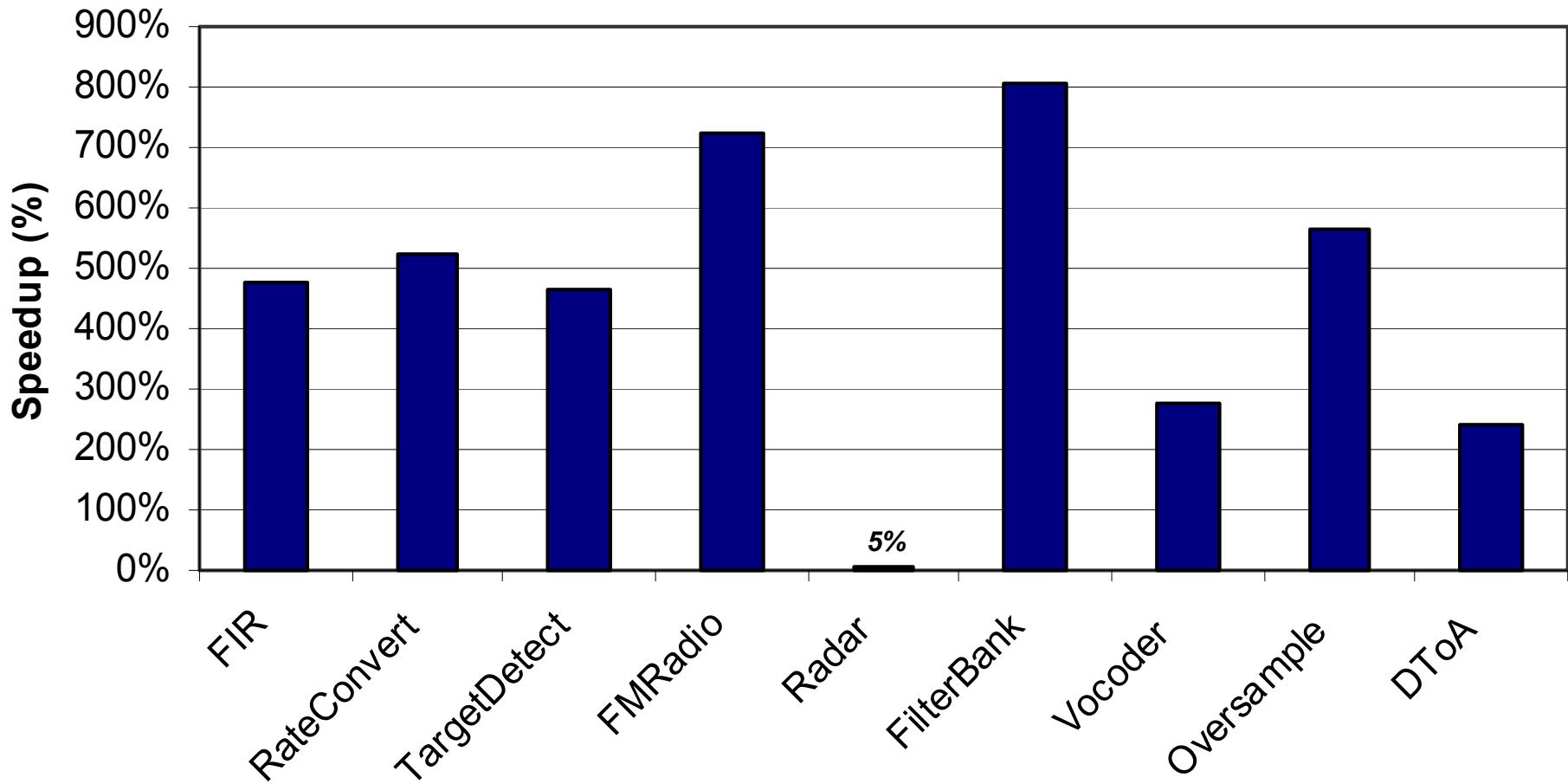
$$C = [32]$$

Filter 2

$$z = y B$$

$$B = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}$$

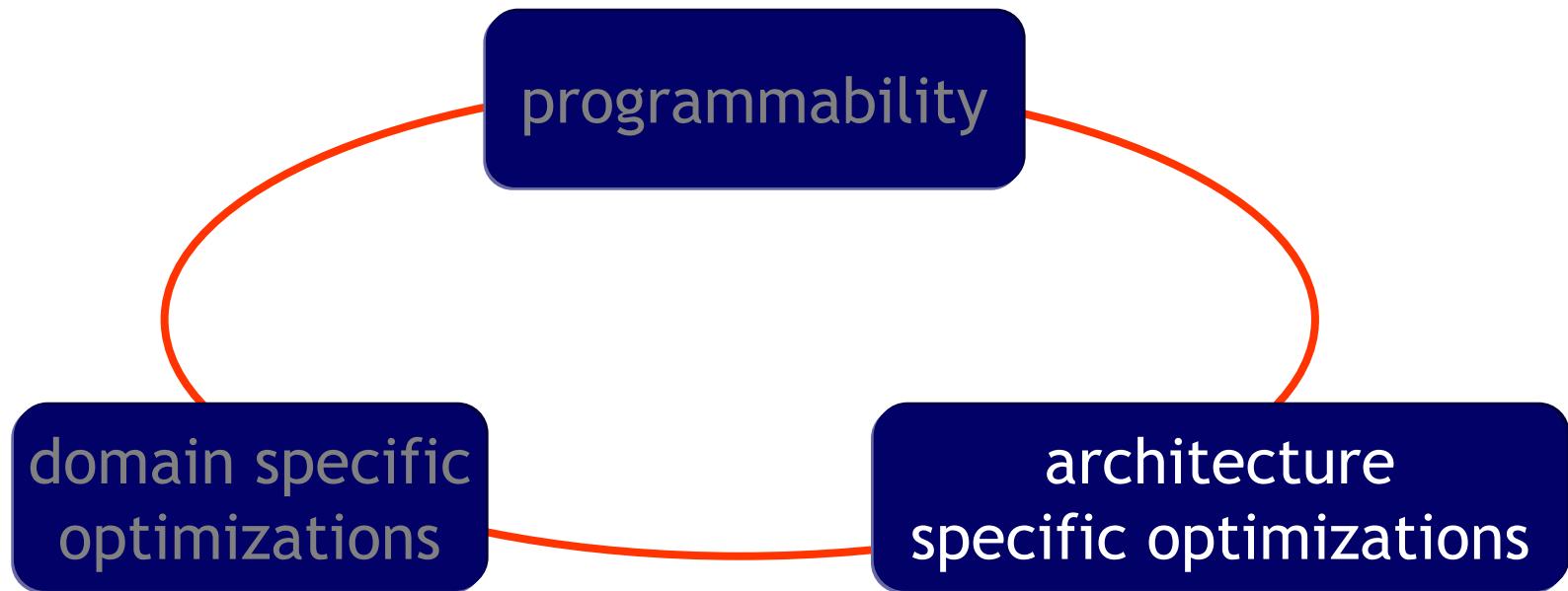
Results from Linear Optimizations



Pentium 4 results compared to baseline StreamIt

The StreamIt Vision

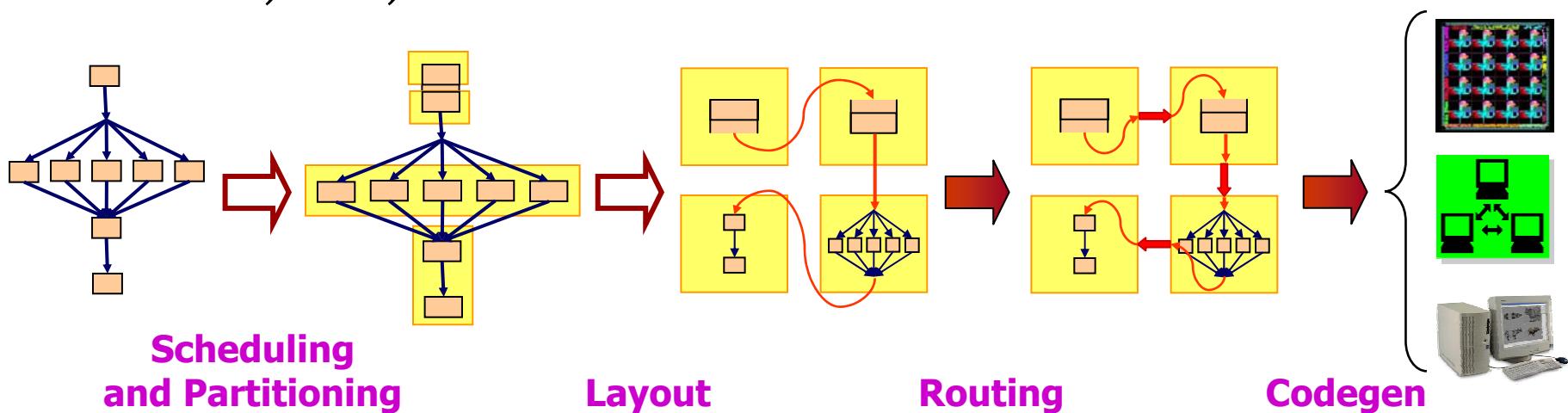
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Core Compilation Technology

- Focused on a common challenges in modern and future architectures
 - MIT Raw fabric architecture
 - Clusters of workstations
 - ARM, x86, and IA-64



- Compiler's role: map the computation and communication pattern to processors, memories, and communication substrates

Compiler Issues

- Load balancing
- Resource utilization
- Fault tolerance
- Dynamic reconfiguration
- ...
- In this talk: cache aware scheduling and partitioning [LCTES '05]

Example Cache Optimization

Baseline

```
for i = 1 to N
    A();
    B();
    C();
end
```

Full Scaling

```
for i = 1 to N
    A();
    for i = 1 to N
        B();
        for i = 1 to N
            C();
```

Cache Aware 1

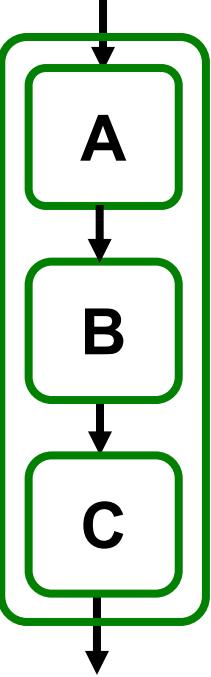
```
for i = 1 to N
    A();
    B();
end
for i = 1 to N
    C();
```

cache size

Working Set Size

	inst	data	inst	data	inst	data
Baseline	A + B + C	A → B → C	A	B	A + B	A → B → C
Full Scaling	A + B + C	A → B → C	A	B	A + B	A → B → C
Cache Aware 1	A + B	A → B → C	A	B	A + B	A → B → C

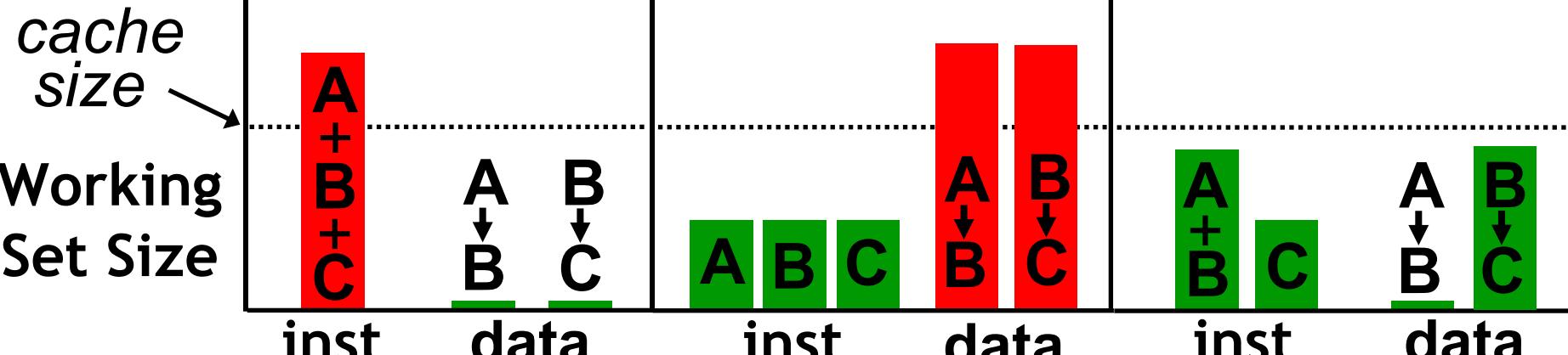
Example Cache Optimization



The diagram illustrates the execution flow from A to C, represented by a vertical sequence of three boxes labeled A, B, and C, with arrows indicating the flow from A down to B, and B down to C.

	Baseline	Full Scaling	Cache Aware 2
	<pre>for i = 1 to N A(); B(); C(); end</pre>	<pre>for i = 1 to N A(); for i = 1 to N B(); for i = 1 to N C();</pre>	<pre>for i = 1 to 64 A(); B(); end for i = 1 to 64 C();</pre>

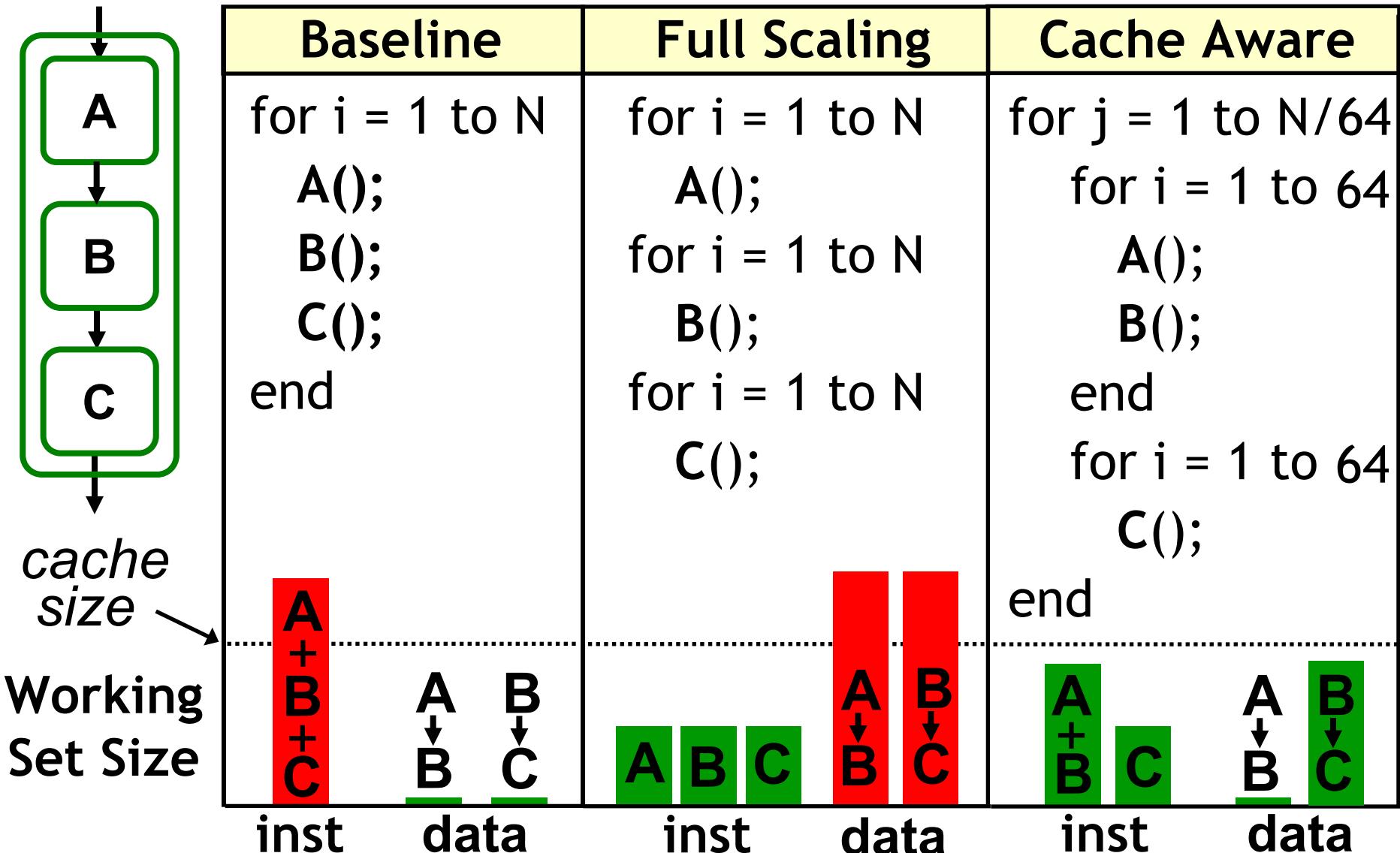
cache size →



The Working Set Size diagram shows the cache utilization over time for three execution models:

- Baseline:** The working set consists of all three instructions (A, B, and C) at all times. It is represented by a red bar labeled "A + B + C" under the "inst" column and a green bar labeled "A ↓ B ↓ C" under the "data" column.
- Full Scaling:** The working set alternates between instruction A (red bar) and instruction B (red bar) for each iteration. It is represented by a red bar labeled "A" under the "inst" column and a green bar labeled "B" under the "data" column.
- Cache Aware 2:** The working set alternates between instruction A (green bar) and instruction C (green bar) for each iteration. It is represented by a green bar labeled "A + B" under the "inst" column and a green bar labeled "C" under the "data" column.

Example Cache Optimization

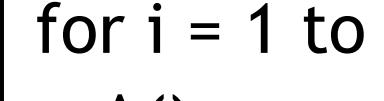


The diagram illustrates the execution flow and cache behavior for three sequential operations (A, B, C) across three stages: Baseline, Full Scaling, and Cache Aware.

Operations:

- Baseline:** for i = 1 to N
 A();
 B();
 C();
end
- Full Scaling:** for i = 1 to N
 A();
 for i = 1 to N
 B();
 C();
- Cache Aware:** for j = 1 to N/64
 for i = 1 to 64
 A();
 B();
 end
 for i = 1 to 64
 C();

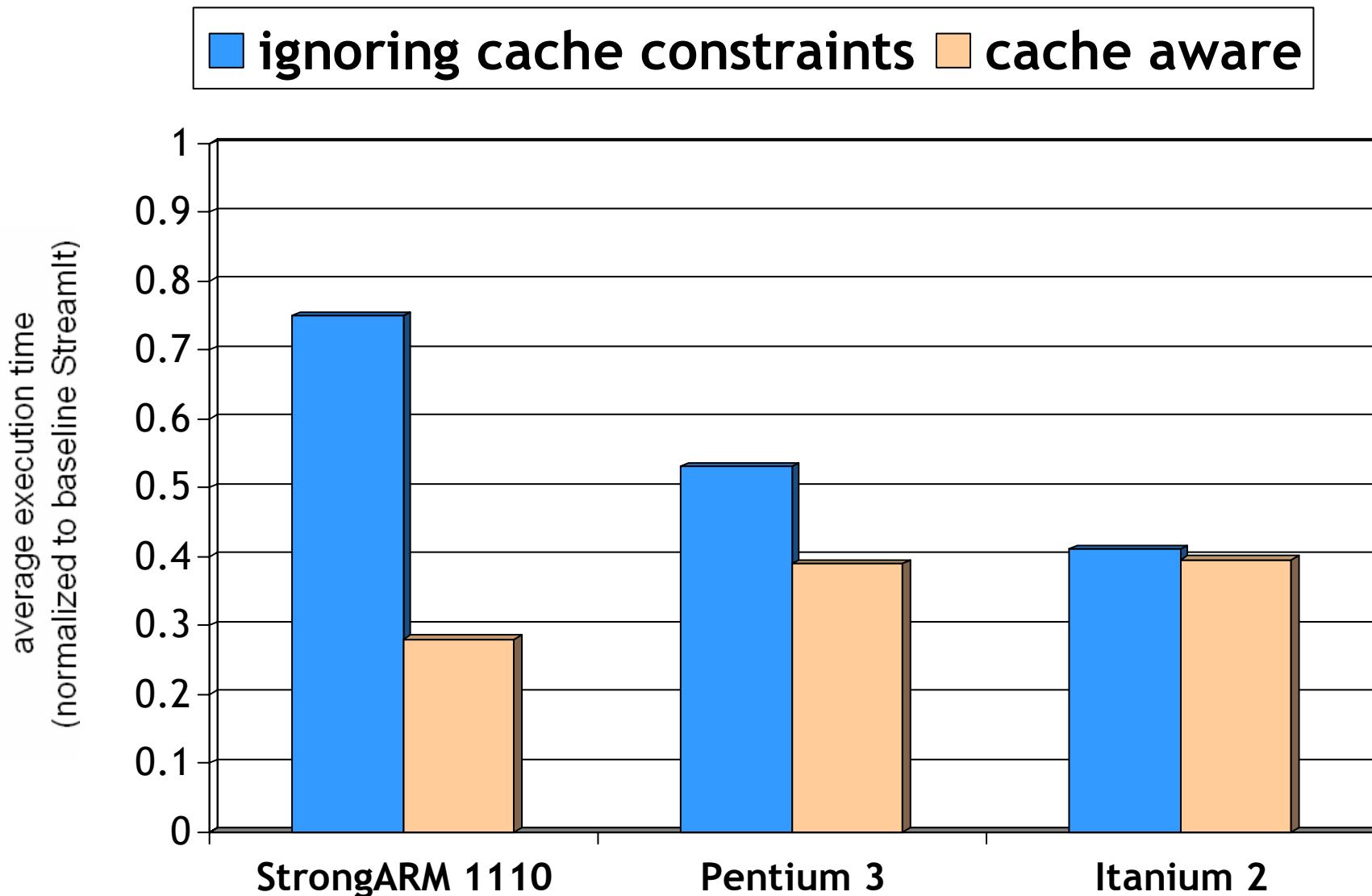
Working Set Size: A red bar indicates the working set size for each stage. In the Cache Aware stage, the working set size is significantly smaller than in the Baseline and Full Scaling stages, demonstrating improved cache utilization.

	Baseline	Full Scaling	Cache Aware
cache size			
Working Set Size	inst:  data: 	inst:  data: 	inst:  data: 
	inst: A data: B C	inst: A data: B C	inst: A data: B C

Evaluation Methodology

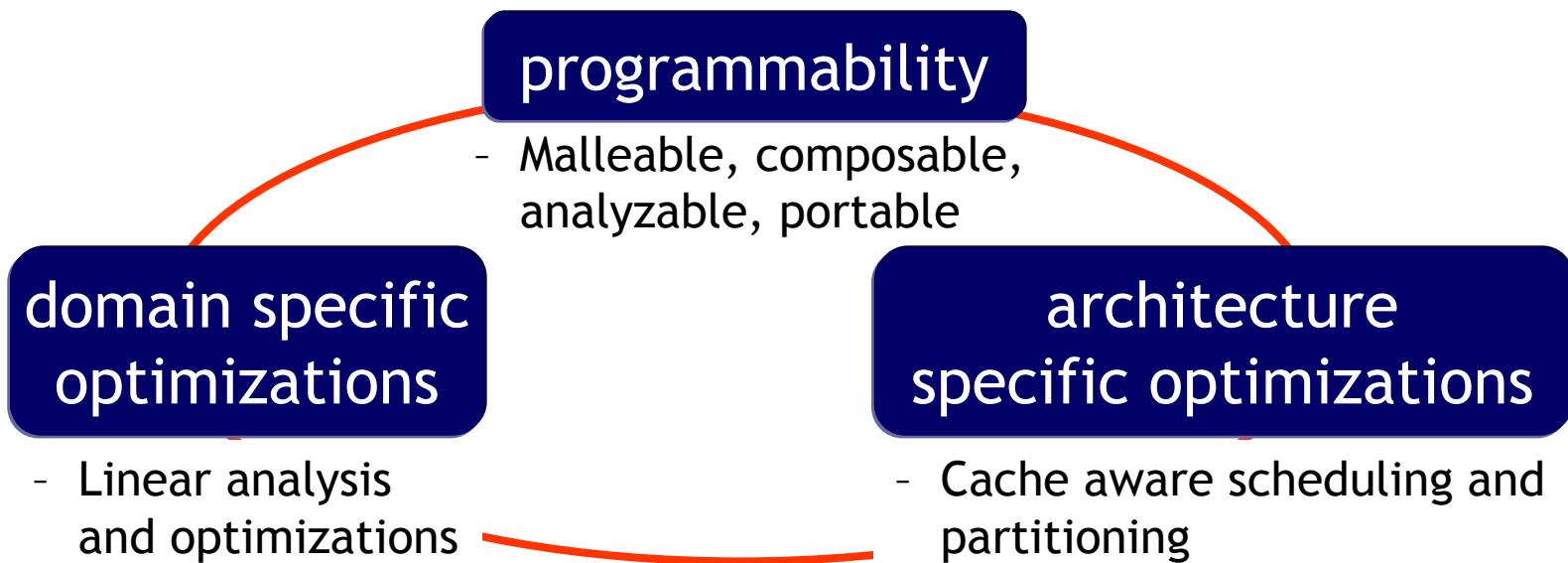
- StreamIt compiler generates C code
 - Baseline StreamIt optimizations
 - Unrolling, constant propagation
 - Compile C code with gcc-v3.4 with -O3 optimizations
- StrongARM 1110 (XScale) embedded processor
 - 370MHz, 16Kb I-Cache, 8Kb D-Cache
 - No L2 Cache (memory 100× slower than cache)
 - Median user time
- Also Pentium 3 and Itanium 2 processors
- Suite of 11 StreamIt Benchmarks

Cache Optimizations Results



Concluding Remarks

- StreamIt improves programmer productivity without compromising performance
 - Easily identify pipeline and data parallelism
 - Expose information for domain specific and architecture specific optimizations



Broader Impact

- Integration into future HPCS languages
 - IBM: X10
- StreamIt for graphics applications
 - Programmable graphics pipeline [Graphics Hardware '05]
- StreamIt for emerging architectures
- Looking for users with interesting applications

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